Graphs in Machine Learning

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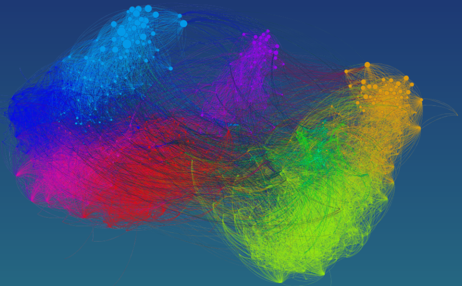
Partially based on material by: Mikhail Belkin, Branislav Kveton



Previous Lecture

- Semi-Supervised Learning
  - Why and when it helps?
  - Self-training
  - Semi-supervised SVMs
- Graph-based semi-supervised learning
- SSL with MinCuts
- Gaussian random fields and harmonic solution
- Regularization of harmonic solution
- Soft-harmonic solution
- Inductive and transductive semi-supervised learning
- Manifold regularization

This Lecture

- Max-Margin Graph Cuts
- Theory of Laplacian-based manifold methods
- Transductive learning stability based bounds
- Online Semi-Supervised Learning
- Online incremental k-centers

Previous Lab Session

- 24. 10. 2016 by Daniele.Calandriello@inria.fr
- Content
  - Graph Construction
  - Test sensitivity to parameters: σ, k, ε
  - Spectral Clustering
  - Spectral Clustering vs. k-means
  - Image Segmentation
- Short written report
- Questions to piazza (without giving away solutions)
- *Deadline:* 7. 11. 2016 Today!
- If you have 32bit OS, send a post to Daniele on piazza

Next Lab Session

- 14. 11. 2016 by Daniele Calandriello
- Content
  - Semi-supervised learning
  - Graph quantization
  - Online face recognizer
- AR: record a video with faces
- Install VM (in case you have not done it yet for TD1)
- Short written report
- Questions to piazza
- *Deadline:* 28. 11. 2016

Final Class projects

- detailed description on the class website
- preferred option: you come up with the topic
- theory/implementation/review or a combination
- one or two people per project (exceptionally three)
- grade 60%: report + short presentation of the **team**
- deadlines
  - 21. 11. 2016 - strongly recommended DL for taking projects
  - 28. 11. 2016 - hard DL for taking projects
  - 05. 01. 2017 - submission of the project report
  - 09. 01. 2017 or later - project presentation
- list of suggested topics on piazza

Master internships in Lille

Master internships in Machine Learning are offered at Inria Lille (Magnet team), **with possibility to pursue a PhD**. Lille is 1h away from Paris, 30min from Brussels, 1h30 from London and 2h30 from Amsterdam.

Available topics:

- **Decentralized machine learning:** Consider a large set of users in a peer-to-peer network, each with a personal dataset. How can users efficiently learn predictive models over the union of their datasets? How can we formally guarantee that their privacy is preserved?
  2 offers: `http://goo.gl/EgYFAs` and `http://goo.gl/EPjSZi`
- **Natural Language Processing:** How to learn dense vector representations of words (or higher-order linguistic entities) from huge text corpora using graph-based spectral algorithms?
  1 offer: `http://goo.gl/0hJjRf`

For more information, contact Aurelien.Bellet@inria.fr

Checkpoint 1

Semi-supervised learning with graphs:

$$\min_{\mathbf{f}\in\{\pm 1\}^{n_l+n_u}} (\infty) \sum_{i=1}^{n_l} w_{ij} \left(f(\mathbf{x}_i) - y_i\right)^2 + \lambda \sum_{i,j=1}^{n_l+n_u} \left(f(\mathbf{x}_i) - f(\mathbf{x}_j)\right)^2$$

Regularized harmonic Solution:

$$\mathbf{f}_u = \left(\mathbf{L}_{uu} + \gamma_g \mathbf{I}\right)^{-1} \left(\mathbf{W}_{ul}\mathbf{f}_l\right)$$

Checkpoint 2

Unconstrained regularization in general:

$$\mathbf{f}^\star = \min_{\mathbf{f}\in\mathbb{R}^N} (\mathbf{f}-\mathbf{y})^\mathsf{T}\mathbf{C}(\mathbf{f}-\mathbf{y}) + \mathbf{f}^\mathsf{T}\mathbf{Q}\mathbf{f}$$

Out of sample extension: Laplacian SVMs

$$f^\star = \arg\min_{f\in\mathcal{H}_\mathcal{K}} \sum_i^{n_l} \max\left(0, 1 - yf\left(\mathbf{x}\right)\right) + \lambda_1 \|f\|_\mathcal{K}^2 + \lambda_2 \mathbf{f}^\mathsf{T}\mathbf{L}\mathbf{f}$$

SSL with Graphs: Laplacian SVMs

$$f^\star = \arg\min_{f \in \mathcal{H}_\mathcal{K}} \sum_{i}^{n_l} \max\left(0, 1 - yf\left(\mathbf{x}\right)\right) + \lambda_1 \|f\|_\mathcal{K}^2 + \lambda_2 \mathbf{f}^\mathsf{T}\mathbf{L}\mathbf{f}$$

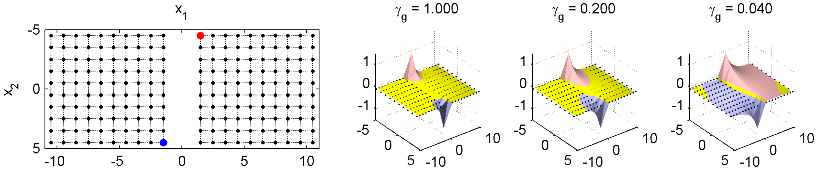
$\mathcal{H}_\mathcal{K}$ is nice and expressive.

Can there be a problem with certain $\mathcal{H}_\mathcal{K}$?

We look for f only in $\mathcal{H}_\mathcal{K}$.

If it is simple (e.g., **linear**) minimization of $\mathbf{f}^\mathsf{T}\mathbf{L}\mathbf{f}$ can perform badly.

Consider again this 2D data and linear $\mathcal{K}$.



SSL with Graphs: Laplacian SVMs

Linear $\mathcal{K} \equiv$ functions with slope α_1 and intercept α_2.

$$\min_{\alpha_1,\alpha_2} \sum_{i}^{n_l} V(f, \mathbf{x}_i, y_i) + \lambda_1 \left[\alpha_1^2 + \alpha_2^2\right] + \lambda_2 \mathbf{f}^\top \mathbf{L} \mathbf{f}$$

For this simple case we can write down $\mathbf{f}^\top\mathbf{L}\mathbf{f}$ explicitly.

$$\begin{aligned}
\mathbf{f}^\top\mathbf{L}\mathbf{f} &= \frac{1}{2}\sum_{i,j} w_{ij}(f(\mathbf{x}_i) - f(\mathbf{x}_j))^2 \\
&= \frac{1}{2}\sum_{i,j} w_{ij}(\alpha_1(\mathbf{x}_{i1} - \mathbf{x}_{j1}) + \alpha_2(\mathbf{x}_{i2} - \mathbf{x}_{j2}))^2 \\
&= \frac{\alpha_1^2}{2}\underbrace{\sum_{i,j} w_{ij}(\mathbf{x}_{i1} - \mathbf{x}_{j1})^2}_{\Delta=218.351} + \frac{\alpha_2^2}{2}\underbrace{\sum_{i,j} w_{ij}(\mathbf{x}_{i2} - \mathbf{x}_{j2})^2}_{\Delta=218.351}
\end{aligned}$$

SSL with Graphs: Laplacian SVMs

2D data and linear $\mathcal{K}$ objective

$$\min_{\alpha_1,\alpha_2} \sum_i^{n_l} V(f,\mathbf{x}_i,y_i) + \left(\lambda_1 + \frac{\lambda_2\Delta}{2}\right)[\alpha_1^2+\alpha_2^2]$$

Setting $\lambda^\star = \left(\lambda_1 + \frac{\gamma_2\Delta}{2}\right)$:

$$\min_{\alpha_1,\alpha_2} \sum_i^{n_l} V(f,\mathbf{x}_i,y_i) + \lambda^\star[\alpha_1^2+\alpha_2^2]$$

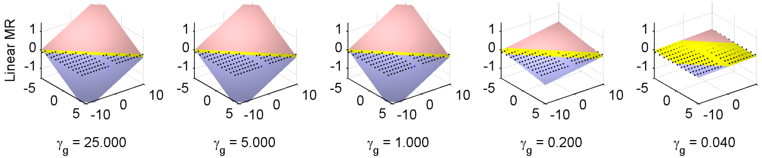
What does this objective function correspond to?

The only influence of unlabeled data is through $\lambda^\star$.

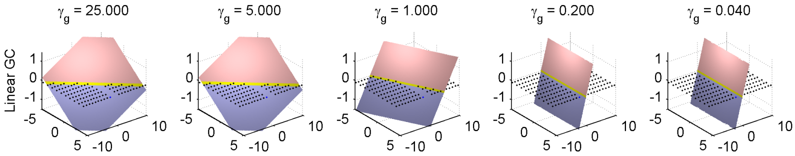
The same value of the objective as for supervised learning for some λ **without the unlabeled data**! This is not good.

SSL with Graphs: Laplacian SVMs

MR for 2D data and linear $\mathcal{K}$ only changes the slope



What would we like to see?



One solution: We use the unlabeled data **before** optimizing over $\mathcal{H}_\mathcal{K}$!

SSL with Graphs: Max-Margin Graph Cuts

Let's take the confident data and use them as true!

$$f^\star = \min_{f\in\mathcal{H}_\mathcal{K}} \sum_{i:|\ell_i^\star|\geq\varepsilon} V(f, \mathbf{x}_i, \mathrm{sgn}(\ell_i^\star)) + \gamma\|f\|_\mathcal{K}^2$$

$$\text{s.t.} \quad \boldsymbol{\ell}^\star = \arg\min_{\boldsymbol{\ell}\in\mathbb{R}^N} \boldsymbol{\ell}^\mathsf{T}(\mathbf{L}+\gamma_g\mathbf{I})\boldsymbol{\ell}$$

$$\text{s.t. } \ell_i = y_i \text{ for all } i = 1,\dots,n_l$$

Wait, but this is what we did not like in self-training!

Will we get into the same trouble?

Representer theorem is still cool:

$$f^\star(\mathbf{x}) = \sum_{i:|f_i^\star|\geq\varepsilon} \alpha_i^\star \mathcal{K}(\mathbf{x}_i, \mathbf{x})$$

SSL with Graphs: Generalization Bounds

Why is this not a witchcraft?

We take GC as an example. MR or HFS are similar.

What kind of guarantees we want?

We may want to bound the **risk**

$$R_P(f) = \mathbb{E}_{P(\mathbf{x})}\left[\mathcal{L}\left(f\left(\mathbf{x}\right), y\left(\mathbf{x}\right)\right)\right]$$

for some **loss**, e.g., 0/1 loss

$$\mathcal{L}(y', y) = \mathbb{1}\{\mathrm{sgn}(y') \neq y\}$$

What makes sense to bound $R_P(f)$ with?

empirical risk + **error terms**

SSL with Graphs: Generalization Bounds

True risk vs. empirical risk

$$R_P(f) = \frac{1}{N}\sum_i (f_i - y_i)^2$$

$$\widehat{R}_P(f) = \frac{1}{n_l}\sum_{i \in l} (f_i - y_i)^2$$

We look for the bound in the form

$$R_P(f) \leq \widehat{R}_P(f) + \text{errors}$$

errors = **transductive** + **inductive**

SSL with Graphs: Generalization Bounds

Bounding inductive error (using classical SLT tools)

With probability $1-\eta$, using Equations 3.15 and 3.24 [**vapnik1995nature**]

$$R_P(f) \leq \frac{1}{n}\sum_i \mathcal{L}(f(\mathbf{x}_i), y_i) + \Delta_I(h, n, \eta).$$

$n \equiv$ number of samples , $h \equiv$ VC dimension of the class

$$\Delta_I(h, n, \eta) = \sqrt{\frac{h(\ln(2n/h)+1) - \ln(\eta/4)}{n}}$$

How to bound $\mathcal{L}(f(\mathbf{x}_i), y_i)$? For any $y_i \in \{-1, 1\}$ and $\ell_i^\star$

$$\mathcal{L}(f(\mathbf{x}_i), y_i) \leq \mathcal{L}(f(\mathbf{x}_i), \mathrm{sgn}(\ell_i^\star)) + (\ell_i^\star - y_i)^2.$$

SSL with Graphs: Generalization Bounds

Bounding transductive error (using stability analysis)

http://www.cs.nyu.edu/~mohri/pub/str.pdf

How to bound $(\ell_i^\star - y_i)^2$?

Bounding $(\ell_i^\star - y_i)^2$ for hard case is difficult $\rightarrow$ we bound soft HFS:

$$\boldsymbol{\ell}^\star = \min_{\boldsymbol{\ell}\in\mathbb{R}^N} (\boldsymbol{\ell} - \mathbf{y})^\mathsf{T}\mathbf{C}(\boldsymbol{\ell} - \mathbf{y}) + \boldsymbol{\ell}^\mathsf{T}\mathbf{Q}\boldsymbol{\ell}$$

Closed form solution

$$\boldsymbol{\ell}^\star = \left(\mathbf{C}^{-1}\mathbf{Q} + \mathbf{I}\right)^{-1}\mathbf{y}$$

SSL with Graphs: Generalization Bounds

Bounding transductive error

$$\boldsymbol{\ell}^\star = \min_{\boldsymbol{\ell}\in\mathbb{R}^N} (\boldsymbol{\ell}-\mathbf{y})^\mathsf{T}\mathbf{C}(\boldsymbol{\ell}-\mathbf{y}) + \boldsymbol{\ell}^\mathsf{T}\mathbf{Q}\boldsymbol{\ell}$$

Think about **stability** of this solution.

Consider two datasets differing in exactly one *labeled* point.

$\mathcal{C}_1 = \mathbf{C}_1^{-1}\mathbf{Q} + \mathbf{I}$ and $\mathcal{C}_2 = \mathbf{C}_2^{-1}\mathbf{Q} + \mathbf{I}$

What is the maximal difference in the solutions?

$$\begin{aligned}
\boldsymbol{\ell}_2^\star - \boldsymbol{\ell}_1^\star &= \mathcal{C}_2^{-1}\mathbf{y}_2 - \mathcal{C}_1^{-1}\mathbf{y}_1 \\
&= \mathcal{C}_2^{-1}(\mathbf{y}_2 - \mathbf{y}_1) - \left(\mathcal{C}_2^{-1} - \mathcal{C}_1^{-1}\right)\mathbf{y}_1 \\
&= \mathcal{C}_2^{-1}(\mathbf{y}_2 - \mathbf{y}_1) - \left(\mathcal{C}_1^{-1}\left[\left(\mathbf{C}_1^{-1} - \mathbf{C}_2^{-1}\right)\mathbf{Q}\right]\mathcal{C}_2^{-1}\right)\mathbf{y}_1
\end{aligned}$$

Note that $\mathbf{v} \in \mathbb{R}^{N\times 1}$, $\lambda_m(A)\|\mathbf{v}\|_2 \le \|A\mathbf{v}\|_2 \le \lambda_M(A)\|\mathbf{v}\|_2$

$$\|\boldsymbol{\ell}_2^\star - \boldsymbol{\ell}_1^\star\|_2 \le \frac{\|\mathbf{y}_2 - \mathbf{y}_1\|_2}{\lambda_m(\mathcal{C}_2)} + \frac{\lambda_M(\mathbf{Q})\|\mathbf{C}_1^{-1} - \mathbf{C}_2^{-1}\|_2 \cdot \|\mathbf{y}_1\|_2}{\lambda_m(\mathcal{C}_2)\lambda_m(\mathcal{C}_1)}$$

SSL with Graphs: Generalization Bounds

Bounding transductive error

$$\boldsymbol{\ell}^\star = \min_{\boldsymbol{\ell}\in\mathbb{R}^N} (\boldsymbol{\ell}-\mathbf{y})^\mathsf{T}\mathbf{C}(\boldsymbol{\ell}-\mathbf{y}) + \boldsymbol{\ell}^\mathsf{T}\mathbf{Q}\boldsymbol{\ell}$$

$$\|\boldsymbol{\ell}_2^\star - \boldsymbol{\ell}_1^\star\|_2 \le \frac{\|\mathbf{y}_2 - \mathbf{y}_1\|_2}{\lambda_m(\mathcal{C}_2)} + \frac{\lambda_M(\mathbf{Q})\|\mathbf{C}_1^{-1} - \mathbf{C}_2^{-1}\|_2 \cdot \|\mathbf{y}_1\|_2}{\lambda_m(\mathcal{C}_2)\lambda_m(\mathcal{C}_1)}$$

Using $\lambda_m(\mathcal{C}) \ge \frac{\lambda_m(\mathbf{Q})}{\lambda_M(\mathbf{C})} + 1$

$$\|\boldsymbol{\ell}_2^\star - \boldsymbol{\ell}_1^\star\|_2 \le \frac{\|\mathbf{y}_2 - \mathbf{y}_1\|_2}{\frac{\lambda_m(\mathbf{Q})}{\lambda_M(\mathbf{C}_1)} + 1} + \frac{\lambda_M(\mathbf{Q})\|\mathbf{C}_1^{-1} - \mathbf{C}_2^{-1}\|_2 \cdot \|\mathbf{y}_1\|_2}{\left(\frac{\lambda_m(\mathbf{Q})}{\lambda_M(\mathbf{C}_2)} + 1\right)\left(\frac{\lambda_m(\mathbf{Q})}{\lambda_M(\mathbf{C}_1)} + 1\right)}$$

SSL with Graphs: Generalization Bounds

Bounding transductive error

$$\|\boldsymbol{\ell}_2^\star - \boldsymbol{\ell}_1^\star\|_\infty \leq \beta \leq \frac{\|\mathbf{y}_2 - \mathbf{y}_1\|_2}{\frac{\lambda_m(\mathbf{Q})}{\lambda_M(\mathbf{C}_1)} + 1} + \frac{\lambda_M(\mathbf{Q})\|\mathbf{C}_1^{-1} - \mathbf{C}_2^{-1}\|_2 \cdot \|\mathbf{y}_1\|_2}{\left(\frac{\lambda_m(\mathbf{Q})}{\lambda_M(\mathbf{C}_2)} + 1\right)\left(\frac{\lambda_m(\mathbf{Q})}{\lambda_M(\mathbf{C}_1)} + 1\right)}$$

Now, let us plug in the values for our problem.

Take $c_l = 1$ and $c_l > c_u$. We have $|y_i| \leq 1$ and $|\ell_i^\star| \leq 1$.

$$\beta \leq 2\left[\frac{\sqrt{2}}{\lambda_m(\mathbf{Q}) + 1} + \sqrt{2n_l}\frac{1 - c_u}{c_u}\frac{\lambda_M(\mathbf{Q})}{(\lambda_m(\mathbf{Q}) + 1)^2}\right]$$

$\mathbf{Q}$ is reg. $\mathbf{L}$: $\lambda_m(\mathbf{Q}) = \lambda_m(\mathbf{L}) + \gamma_g$ and $\lambda_M(\mathbf{Q}) = \lambda_M(\mathbf{L}) + \gamma_g$

$$\beta \;\leq\; 2\left[\frac{\sqrt{2}}{\gamma_g + 1} + \sqrt{2n_l}\frac{1 - c_u}{c_u}\frac{\lambda_M(\mathbf{L}) + \gamma_g}{\gamma_g^2 + 1}\right]$$

This algorithm is β-stable!

SSL with Graphs: Generalization Bounds

Bounding transductive error

http://web.cse.ohio-state.edu/~mbelkin/papers/RSS_COLT_04.pdf

By the generalization bound of Belkin [**belkin2004regularization**]

$$R_P^W(\boldsymbol{\ell}^\star) \;\leq\; \widehat{R}_P^W(\boldsymbol{\ell}^\star) + \underbrace{\beta + \sqrt{\frac{2\ln(2/\delta)}{n_l}}(n_l\beta + 4)}_{\text{transductive error } \Delta_T(\beta, n_l, \delta)}$$

$$\beta \;\leq\; 2\left[\frac{\sqrt{2}}{\gamma_g + 1} + \sqrt{2n_l}\frac{1 - c_u}{c_u}\frac{\lambda_M(\mathbf{L}) + \gamma_g}{\gamma_g^2 + 1}\right]$$

holds with probability $1 - \delta$, where

$$R_P^W(\boldsymbol{\ell}^\star) \;=\; \frac{1}{N}\sum_i (\ell_i^\star - y_i)^2$$

$$\widehat{R}_P^W(\boldsymbol{\ell}^\star) \;=\; \frac{1}{n_l}\sum_{i \in l} (\ell_i^\star - y_i)^2.$$

SSL with Graphs: Generalization Bounds

Bounding transductive error

$$R_P^W(\ell^\star) \;\leq\; \widehat{R}_P^W(\ell^\star) + \underbrace{\beta + \sqrt{\frac{2\ln(2/\delta)}{n_l}}(n_l\beta + 4)}_{\text{transductive error } \Delta_T(\beta, n_l, \delta)}$$

$$\beta \;\leq\; 2\left[\frac{\sqrt{2}}{\gamma_g + 1} + \sqrt{2n_l}\frac{1 - c_u}{c_u}\frac{\lambda_M(\mathbf{L}) + \gamma_g}{\gamma_g^2 + 1}\right]$$

Does the bound say anything useful?

1) The error is controlled.

2) Practical when error $\Delta_T(\beta, n_l, \delta)$ decreases at rate $O(n_l^{-\frac{1}{2}})$.
Achieved when $\beta = O(1/n_l)$. That is, $\gamma_g = \Omega(n_l^{\frac{3}{2}})$.

We have an idea how to set γ_g!

SSL with Graphs: Generalization Bounds

Combining inductive + transductive error

With probability $1-(\eta+\delta)$.

$$R_P(f) \leq \frac{1}{n}\sum_i \mathcal{L}(f(\mathbf{x}_i), \mathrm{sgn}(\ell_i^\star)) + \widehat{R}_P^W(\boldsymbol{\ell}^\star) + \Delta_T(\beta, n_l, \delta) + \Delta_I(h, N, \eta)$$

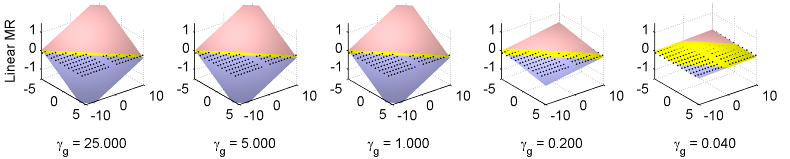
We need to account for ε. With probability $1-(\eta+\delta)$.

$$R_P(f) \leq \frac{1}{n}\sum_{i:|\ell_i^\star|\geq\varepsilon} \mathcal{L}(f(\mathbf{x}_i), \mathrm{sgn}(\ell_i^\star)) + \frac{2\varepsilon n_\varepsilon}{N} + \widehat{R}_P^W(\boldsymbol{\ell}^\star) + \Delta_T(\beta, n_l, \delta) + \Delta_I(h, N, \eta)$$

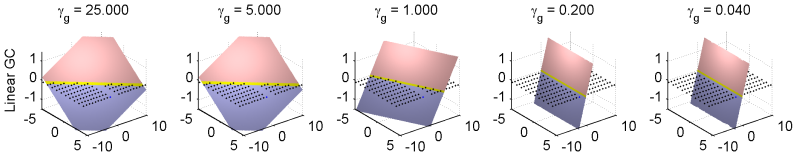
We should have $\varepsilon \leq n_l^{-1/2}$!

SSL with Graphs: LapSVMs and MM Graph Cuts

MR for 2D data and **linear** $\mathcal{K}$ only changes the slope

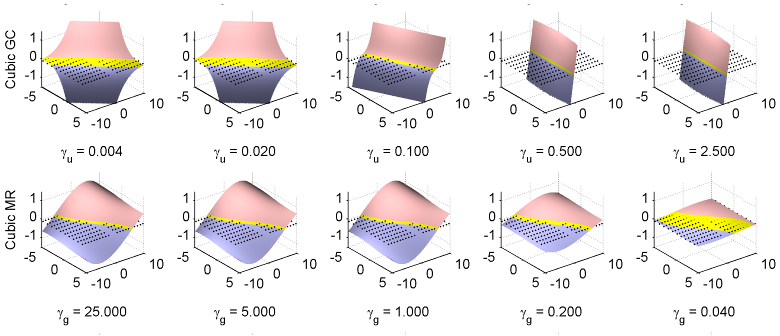


MMGC for 2D data and **linear** $\mathcal{K}$ works as we want



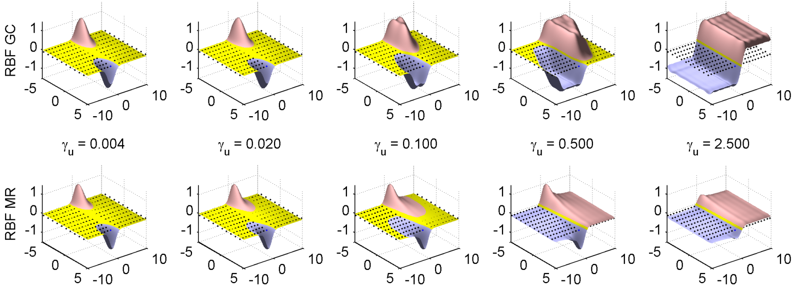
SSL with Graphs: LapSVMs and MM Graph Cuts

MR for 2D data and **cubic** $\mathcal{K}$ is also not so good

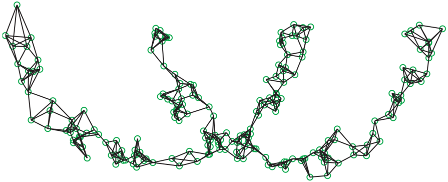


SSL with Graphs: LapSVMs and MM Graph Cuts

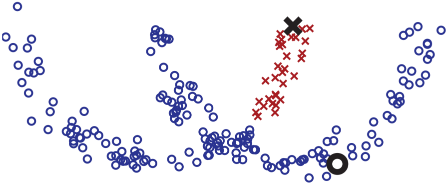
MMGC and MR for 2D data and RBF $\mathcal{K}$



SSL with Graphs



(a) 4-NN graph



(b) Harmonic function predictions

Graph-based SSL is obviously sensitive to graph construction!

Online SSL with Graphs

Offline learning setup

Given $\{\mathbf{x}_i\}_{i=1}^{N}$ from $\mathbb{R}^d$ and $\{y_i\}_{i=1}^{n_l}$, with $n_l \ll n$, find $\{y_i\}_{i=n_l+1}^{N}$ (**transductive**) or find f predicting y well beyond that (**inductive**).

Online learning setup

At the beginning: $\{\mathbf{x}_i, y_i\}_{i=1}^{n_l}$ from $\mathbb{R}^d$

At time t:

- receive $\mathbf{x}_t$
- predict y_t

Online SSL with Graphs

Online HFS: Straightforward solution

1: **while** new unlabeled example $\mathbf{x}_t$ comes **do**
2: Add $\mathbf{x}_t$ to graph $G(\mathbf{W})$
3: Update $\mathbf{L}_t$
4: Infer labels

$$\mathbf{f}_u = \left(\mathbf{L}_{uu} + \gamma_g \mathbf{I}\right)^{-1}\left(\mathbf{W}_{ul}\mathbf{f}_l\right)$$

5: Predict $\widehat{y}_t = \operatorname{sgn}\left(\mathbf{f}_u\left(t\right)\right)$
6: **end while**

What is wrong with this solution?

The cost and memory of the operations.

What can we do?

Online SSL with Graphs

Let's keep only k vertices!

Limit memory to k **centroids** with $\widetilde{\mathbf{W}}^{\mathrm{q}}$ weights.

Each centroids represents *several* others.

Diagonal $\mathbf{V} \equiv$ **multiplicity**. We have $\mathbf{V}_{ii}$ copies of centroid i.

Can we compute it compactly? Compact harmonic solution.

$$\boldsymbol{\ell}^{\mathrm{q}} = (\mathbf{L}^{\mathrm{q}}_{uu} + \gamma_g V)^{-1}\mathbf{W}^{\mathrm{q}}_{ul}\boldsymbol{\ell}_l \quad \text{where} \quad \mathbf{W}^{\mathrm{q}} = V\widetilde{\mathbf{W}}^{\mathrm{q}}V$$

Proof? Using electric circuits.

Why do we keep the multiplicities?

Online SSL with Graphs

Online HFS with Graph Quantization

Input
  k number of representative nodes
Initialization
  $\mathbf{V}$ matrix of multiplicities with 1 on diagonal
while new unlabeled example $\mathbf{x}_t$ comes **do**
  Add $\mathbf{x}_t$ to graph G
  if $\#$ nodes $> k$ **then**
    quantize G
  end if
  Update $\mathbf{L}_t$ of $G(\mathbf{VWV})$
  Infer labels
  Predict $\widehat{y}_t = \operatorname{sgn}(\mathbf{f}_u(t))$
end while

Online SSL with Graphs: Graph Quantization

An idea: incremental k-centers

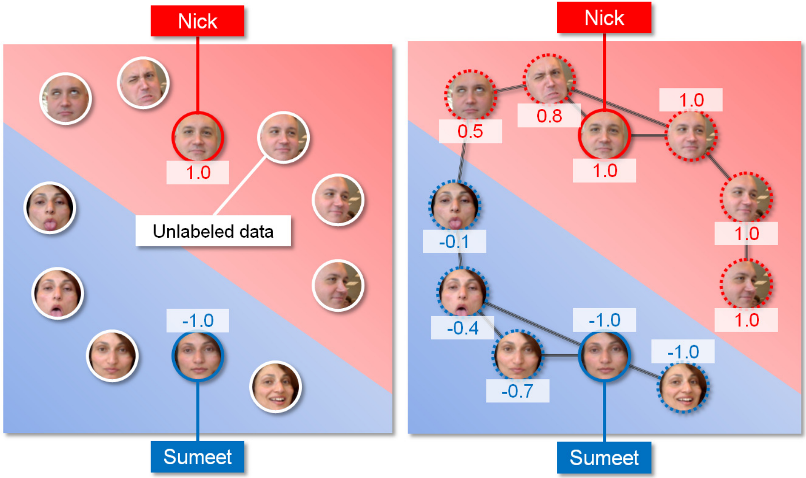
Doubling algorithm of Charikar et al. [**charikar1997incremental**]

Keeps up to k centers $C_t = \{\mathbf{c}_1, \mathbf{c}_2, \dots\}$ with

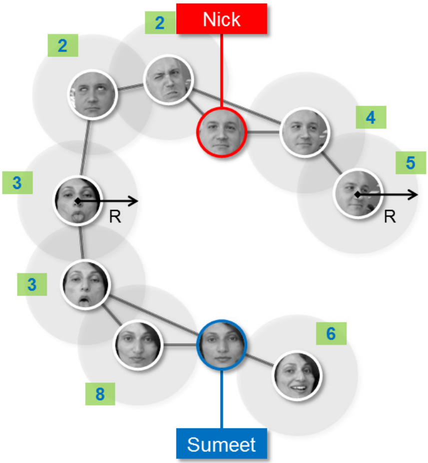
- Distance $\mathbf{c}_i, \mathbf{c}_j \in C_t$ is at least $\geq R$
- For each new $\mathbf{x}_t$, distance to some $\mathbf{c}_i \in C_t$ is less than R.
- $|C_t| \leq k$
- if not possible, R is doubled

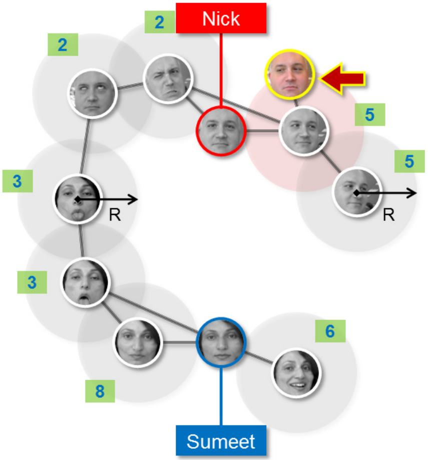
Online SSL with Graphs: Graph Quantization

Online SSL with Graphs: Graph Quantization

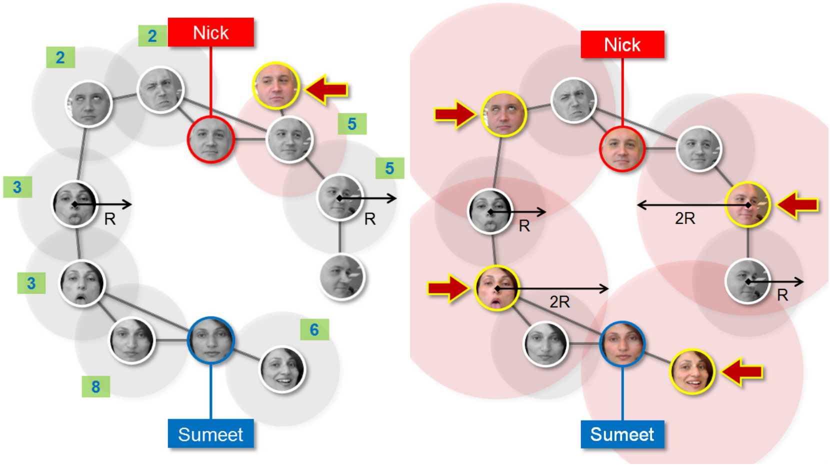


Online SSL with Graphs: Graph Quantization





Online SSL with Graphs: Graph Quantization



Online SSL with Graphs: Graph Quantization

Doubling algorithm [**charikar1997incremental**]

To reduce growth of R, we use $R \leftarrow m \times R$, with $m \geq 1$

C_t is changing. How far can $\mathbf{x}$ be from some $\mathbf{c}$?

$$R + \frac{R}{m} + \frac{R}{m^2} + \cdots = R\left(1 + \frac{1}{m} + \frac{1}{m^2} + \cdots\right) = \frac{Rm}{m-1}$$

Guarantees: $(1+\varepsilon)$-approximation algorithm.

Why not incremental k-means?

Online SSL with Graphs: Graph Quantization

Online k-centers

an unlabeled $\mathbf{x}_t$, a set of centroids C_{t-1}, multiplicities $\mathbf{v}_{t-1}$
if $(|C_{t-1}| = k+1)$ **then**
  $R \leftarrow mR$
  greedily repartition C_{t-1} into C_t such that:
    no two vertices in C_t are closer than R
    for any $\mathbf{c}_i \in C_{t-1}$ exists $\mathbf{c}_j \in C_t$ such that $d(\mathbf{c}_i, \mathbf{c}_j) < R$
  update $\mathbf{v}_t$ to reflect the new partitioning
else
  $C_t \leftarrow C_{t-1}$
  $\mathbf{v}_t \leftarrow \mathbf{v}_{t-1}$
end if
if $\mathbf{x}_t$ is closer than R to any $\mathbf{c}_i \in C_t$ **then**
  $\mathbf{v}_t(i) \leftarrow \mathbf{v}_t(i) + 1$
else
  $\mathbf{v}_t(|C_t| + 1) \leftarrow 1$
end if

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