

# **Graphs in Machine Learning**

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Partially based on material by: Tomáš Kocák

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MVA 2015/2016

#### Last Lecture

- Examples of applications of online SSL
- Analysis of online SSL
- SSL Learnability
- When does graph-based SSL provably help?
- Scaling harmonic functions to millions of samples



#### **Previous Lab Session**

- ▶ 16. 11. 2015 by Daniele Calandriello
- Content
  - Semi-supervised learning
  - Graph quantization
  - Online face recognizer
- Short written report
- Questions to piazza
- Deadline: 30. 11. 2015
- http://researchers.lille.inria.fr/~calandri/teaching.html



#### **This Lecture**

- Online decision-making on graphs
- Graph bandits
- Smoothness of rewards (preferences) on a given graph
- Observability graphs
- Exploiting side information

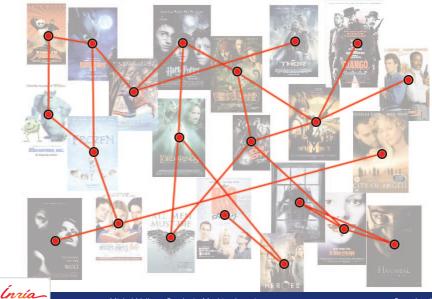


### **Final Class projects**

- detailed description on the class website
- preferred option: you come up with the topic
- theory/implementation/review or a combination
- one or two people per project (exceptionally three)
- ▶ grade 60%: report + short presentation of the team
- deadlines
  - ▶ 23. 11. 2015 strongly recommended DL for taking projects
  - ▶ 30. 11. 2015 hard DL for taking projects
  - ▶ 06. 01. 2015 submission of the project report
  - 11. 01. 2016 (or later) project presentation
- list of suggested topics on piazza



#### **Online Decision Making on Graphs**



### **Online Decision Making on Graphs: Smoothness**

Sequential decision making in structured settings

- we are asked to pick a node (or a few nodes) in a graph
- the graph encodes some structural property of the setting
- goal: maximize the sum of the outcomes
- application: recommender systems
- First application: Exploiting smoothness
  - fixed graph
  - iid outcomes
  - neighboring nodes have similar outcomes

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# **Online Decision Making on Graphs**

Movie recommendation: (in each time step)

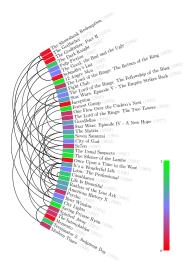
- Recommend movies to a **single user**.
- Good prediction after a few steps ( $T \ll N$ ).

#### Goal:

Maximize overall reward (sum of ratings).

#### Assumptions:

- Unknown reward function  $f: V(G) \rightarrow \mathbb{R}$ .
- Function *f* is **smooth** on a graph.
- Neighboring movies  $\Rightarrow$  similar preferences.
- ► Similar preferences ⇒ neighboring movies.





#### Let's be lazy: Ignore the structure!



This is an multi-armed bandit problem!

The performance depends on the number of movies (N arms).

Worst case regret (to the best fixed strategy)  $R_T = \mathcal{O}\left(\sqrt{NT}\right)$ 

What is *N* for imdb.com? 3,538,545 http://www.imdb.com/stats



#### Let's be lazy: Ignore the structure!

Another problem of the typical bandits strategies for recommendation?

If there is no information shared, we need to try all of the options!

UCB/MOSS and likely TS start with pulling each of the arms once

This is a problem both algorithmically and theoretically ....

Watch all the movies and then I tell you which one you like ....

What do we need for movie recommendation?

An algorithm useful in the case  $T \ll N!$ 

Exploiting the structure is a must!



#### **Recap: Smooth graph functions**

- $\mathbf{f} = (f_1, \dots, f_N)^{\mathsf{T}}$ : Vector of function values.
- Let  $\mathbf{L} = \mathbf{Q} \mathbf{\Lambda} \mathbf{Q}^{\mathsf{T}}$  be the eigendecomposition of the Laplacian.
  - Diagonal matrix Λ whose diagonal entries are eigenvalues of L.
  - Columns of Q are eigenvectors of L.
  - Columns of **Q** form a basis.

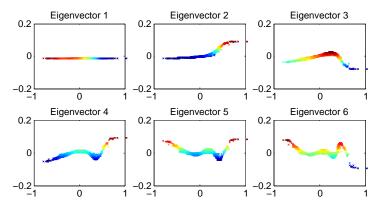
• 
$$\alpha^*$$
: Unique vector such that  $\mathbf{Q}\alpha^* = \mathbf{f}$  Note:  $\mathbf{Q}^{\mathsf{T}}\mathbf{f} = \alpha^*$ 

$$S_G(\mathbf{f}) = \mathbf{f}^{\mathsf{T}} \mathbf{L} \mathbf{f} = \mathbf{f}^{\mathsf{T}} \mathbf{Q} \mathbf{\Lambda} \mathbf{Q}^{\mathsf{T}} \mathbf{f} = \boldsymbol{\alpha}^{*\mathsf{T}} \mathbf{\Lambda} \boldsymbol{\alpha}^* = \| \boldsymbol{\alpha}^* \|_{\mathbf{\Lambda}}^2 = \sum_{i=1}^N \lambda_i (\alpha_i^*)^2$$

Smoothness and <u>regularization</u>: Small value of (a)  $S_G(\mathbf{f})$  (b)  $\Lambda$  norm of  $\alpha^*$  (c)  $\alpha_i^*$  for large  $\lambda_i$ 



# Smooth graph functions: Flixster eigenvectors



Eigenvectors from the Flixster data corresponding to the smallest few eigenvalues of the graph Laplacian projected onto the first principal component of data. Colors indicate the values.



#### **Online Learning Setting - Bandit Problem**

Learning setting for a bandit algorithm  $\boldsymbol{\pi}$ 

- In each time t step choose a node  $\pi(t)$ .
- ▶ the  $\pi(t)$ -th row  $\mathbf{x}_{\pi(t)}$  of the matrix **Q** corresponds to the arm  $\pi(t)$ .
- ► Obtain noisy reward  $r_t = \mathbf{x}_{\pi(t)}^{\mathsf{T}} \boldsymbol{\alpha}^* + \varepsilon_t$ . Note:  $\mathbf{x}_{\pi(t)}^{\mathsf{T}} \boldsymbol{\alpha}^* = f_{\pi(t)}$

►  $\varepsilon_t$  is *R*-sub-Gaussian noise.  $\forall \xi \in \mathbb{R}, \mathbb{E}[e^{\xi \varepsilon_t}] \leq \exp(\xi^2 R^2/2)$ 

Minimize cumulative regret

$$R_T = T \max_a (\mathbf{x}_a^{\mathsf{T}} \boldsymbol{lpha}^*) - \sum_{t=1}^T \mathbf{x}_{\pi(t)}^{\mathsf{T}} \boldsymbol{lpha}^*.$$

What is a good result?

Can't we just use *linear bandits*?



# **Online Decision Making on Graphs: Smoothness**

- Linear bandit algorithms
  - ► LinUCB
    - Regret bound  $\approx D\sqrt{T \ln T}$
  - ► LinearTS
    - Regret bound  $\approx D\sqrt{T \ln N}$

(Li et al., 2010)

(Agrawal and Goyal, 2013)

**Note:** *D* is ambient dimension, in our case *N*, length of  $x_i$ . Number of actions, e.g., all possible movies  $\rightarrow$  **HUGE**!

Spectral bandit algorithms

- SpectralUCB
  - Regret bound  $\approx d\sqrt{T \ln T}$
  - Operations per step: D<sup>2</sup>N
- SpectralTS
  - Regret bound  $\approx d\sqrt{T \ln N}$
  - Operations per step:  $D^2 + DN$

**Note:** *d* is effective dimension, usually much smaller than *D*.



SequeL - 14/40

(Valko et al., ICML 2014)

(Kocák et al., AAAI 2014)

#### **Effective dimension**

**Effective dimension:** Largest *d* such that

$$(d-1)\lambda_d \leq rac{\mathcal{T}}{\log(1+\mathcal{T}/\lambda)}.$$

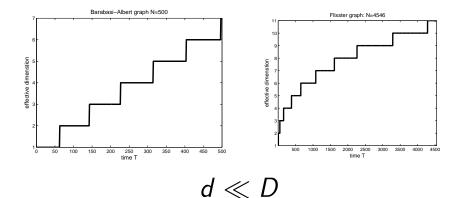
- Function of time horizon and graph properties
- $\lambda_i$ : *i*-th smallest eigenvalue of **A**.
- $\lambda$ : Regularization parameter of the algorithm.

#### **Properties:**

- *d* is small when the coefficients  $\lambda_i$  grow rapidly above time.
- ► *d* is related to the number of "non-negligible" dimensions.
- ▶ Usually *d* is much smaller than *D* in real world graphs.
- Can be computed beforehand.



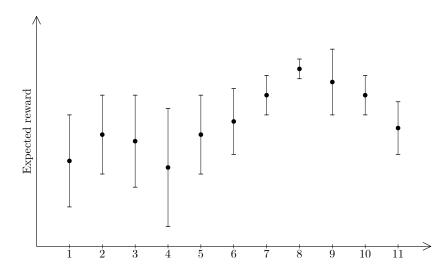
#### Effective dimension vs. Ambient dimension



Note: In our setting T < N = D.



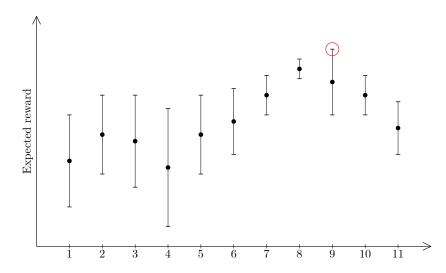
# UCB-style algorithms: Estimate





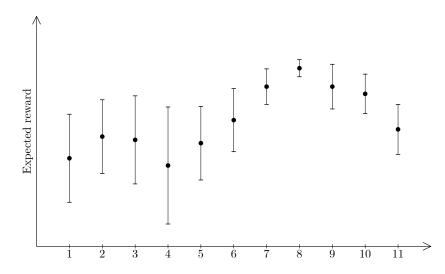
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# UCB-style algorithms: Sample





#### UCB-style algorithms: Estimate ...





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## **SpectralUCB**

Given a vector of weights  $\alpha$ , we define its  $\mathbf{\Lambda}$  norm as

$$\|\boldsymbol{\alpha}\|_{\boldsymbol{\Lambda}} = \sqrt{\sum_{k=1}^{N} \lambda_k \alpha_k^2} = \sqrt{\boldsymbol{\alpha}^{\mathsf{T}} \boldsymbol{\Lambda} \boldsymbol{\alpha}},$$

and fit the ratings  $r_v$  with a (regularized) least-squares estimate

$$\widehat{\boldsymbol{\alpha}}_t = \operatorname*{arg\,min}_{\boldsymbol{\alpha}} \left( \sum_{\nu=1}^t \left[ \langle \boldsymbol{\mathsf{x}}_{\nu}, \boldsymbol{\alpha} \rangle - r_{\nu} \right]^2 + \|\boldsymbol{\alpha}\|_{\boldsymbol{\Lambda}}^2 \right).$$

 $\|\alpha\|_{\Lambda}$  is a penalty for non-smooth combinations of eigenvectors.

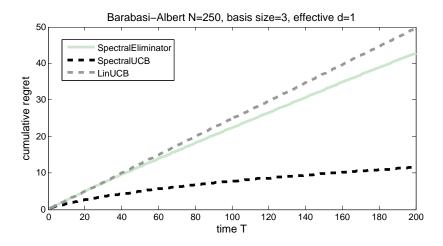


# **SpectralUCB**

1: Input: 2:  $N, T, \{\Lambda_L, \mathbf{Q}\}, \lambda, \delta, R, C$ 3: Run: 4:  $\Lambda \leftarrow \Lambda_1 + \lambda_1$ 5:  $d \leftarrow \max\{d: (d-1)\lambda_d \leq T/\ln(1+T/\lambda)\}$ 6: for t = 1 to T do 7: Update the basis coefficients  $\hat{\alpha}$ :  $\mathbf{X}_t \leftarrow [\mathbf{x}_{\pi(1)}, \dots, \mathbf{x}_{\pi(t-1)}]^{\mathsf{T}}$ 8: 9:  $\mathbf{r} \leftarrow [r_1, \ldots, r_{t-1}]^{\mathsf{T}}$ 10:  $\mathbf{V}_t \leftarrow \mathbf{X}_t \mathbf{X}_t^{\mathsf{T}} + \mathbf{\Lambda}$ 11:  $\widehat{\boldsymbol{\alpha}}_t \leftarrow \mathbf{V}_t^{-1} \mathbf{X}_t^{\mathsf{T}} \mathbf{r}$ 12:  $c_t \leftarrow 2R\sqrt{d\ln(1+t/\lambda)+2\ln(1/\delta)}+C$  $\pi(t) \leftarrow \arg \max_{a} \left( \mathbf{x}_{a}^{\mathsf{T}} \widehat{\alpha} + c_{t} \| \mathbf{x}_{a} \|_{\mathbf{V}_{t}^{-1}} \right)$ 13: 14: Observe the reward  $r_{t}$ 15: end for

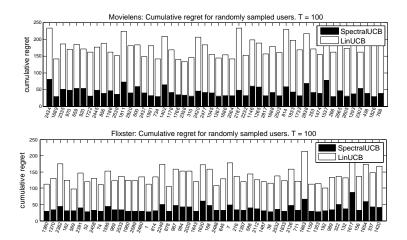


### **SpectralUCB: Synthetic experiment**





#### SpectralUCB: Movie data experiments





- ► *d*: Effective dimension.
- $\lambda$ : Minimal eigenvalue of  $\mathbf{\Lambda} = \mathbf{\Lambda}_{\mathbf{L}} + \lambda \mathbf{I}$ .
- C: Smoothness upper bound,  $\|\alpha^*\|_{\mathbf{A}} \leq C$ .

• 
$$\mathbf{x}_i^{\mathsf{T}} \boldsymbol{\alpha}^* \in [-1, 1]$$
 for all *i*.

The **cumulative regret**  $R_T$  of **SpectralUCB** is with probability  $1 - \delta$  bounded as

$$R_{T} \leq \left(8R\sqrt{d\ln\frac{\lambda+T}{\lambda}+2\ln\frac{1}{\delta}}+4C+4\right)\sqrt{dT\ln\frac{\lambda+T}{\lambda}}.$$

$$R_T \approx d\sqrt{T \ln T}$$



- Derivation of the confidence ellipsoid for  $\hat{\alpha}$  with probability  $1 \delta$ .
  - Using analysis of OFUL (Abbasi-Yadkori et al., 2011)

$$|\mathbf{x}^{\mathsf{T}}(\widehat{\boldsymbol{\alpha}} - \boldsymbol{\alpha}^*)| \leq \|\mathbf{x}\|_{\mathbf{V}_t^{-1}} \left( R_{\sqrt{2 \ln \left(\frac{|\mathbf{V}_t|^{1/2}}{\delta |\mathbf{A}|^{1/2}}\right)} + C \right)$$

- Regret in one time step:  $r_t = \mathbf{x}_*^{\mathsf{T}} \boldsymbol{\alpha}^* \mathbf{x}_{\pi}^{\mathsf{T}} t \boldsymbol{\alpha}^* \leq 2c_t \|\mathbf{x}_{\pi(t)}\|_{\mathbf{V}_t^{-1}}$

$$R_{T} = \sum_{t=1}^{T} r_{t} \leq \sqrt{T \sum_{t=1}^{T} r_{t}^{2}} \leq 2(\frac{V}{C_{T}} + 1) \sqrt{2T \ln \frac{|V_{T}|}{|\Lambda|}}$$

• Upperbound for  $\ln(|\mathbf{V}_t|/|\mathbf{\Lambda}|)$ 

$$\ln \frac{|\mathbf{V}_t|}{|\mathbf{A}|} \leq \ln \frac{|\mathbf{V}_{\mathcal{T}}|}{|\mathbf{A}|} \leq 2d \ln \left(\frac{\lambda + \mathcal{T}}{\lambda}\right)$$



Sylvester's determinant theorem:

$$|\mathbf{A} + \mathbf{x}\mathbf{x}^{\mathsf{T}}| = |\mathbf{A}||\mathbf{I} + \mathbf{A}^{-1}\mathbf{x}\mathbf{x}^{\mathsf{T}}| = |\mathbf{A}|(1 + \mathbf{x}^{\mathsf{T}}\mathbf{A}^{-1}\mathbf{x})$$

Goal:

- $\blacktriangleright$  Upperbound determinant  $|\mathbf{A} + \mathbf{x}\mathbf{x}^{\scriptscriptstyle\mathsf{T}}|$  for  $\|\mathbf{x}\|_2 \leq 1$
- ▶ Upperbound x<sup>T</sup>A<sup>-1</sup>x

$$\mathbf{x}^{\mathsf{T}}\mathbf{A}^{-1}\mathbf{x} = \mathbf{x}^{\mathsf{T}}\mathbf{Q}\mathbf{\Lambda}^{-1}\mathbf{Q}^{\mathsf{T}}\mathbf{x} = \mathbf{y}^{\mathsf{T}}\mathbf{\Lambda}^{-1}\mathbf{y} = \sum_{i=1}^{N} \lambda_{i}^{-1}y_{i}^{2}$$

▶  $\|\mathbf{y}\|_2 \le 1.$ 

- **y** is a canonical vector.
- $\mathbf{x} = \mathbf{Q}\mathbf{y}$  is an eigenvector of  $\mathbf{A}$ .



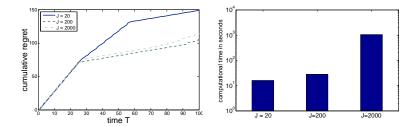
**Corollary**: Determinant  $|\mathbf{V}_{\mathcal{T}}|$  of  $\mathbf{V}_{\mathcal{T}} = \mathbf{\Lambda} + \sum_{t=1}^{\mathcal{T}} \mathbf{x}_t \mathbf{x}_t^{\mathsf{T}}$  is maximized when all  $\mathbf{x}_t$  are aligned with axes.

$$\begin{split} |\mathbf{V}_{\mathcal{T}}| &\leq \max_{\sum t_i = \mathcal{T}} \prod (\lambda_i + t_i) \\ \ln \frac{|\mathbf{V}_{\mathcal{T}}|}{|\mathbf{\Lambda}|} &\leq \max_{\sum t_i = \mathcal{T}} \sum \ln \left(1 + \frac{t_i}{\lambda_i}\right) \\ \ln \frac{|\mathbf{V}_{\mathcal{T}}|}{|\mathbf{\Lambda}|} &\leq \sum_{i=1}^{d} \ln \left(1 + \frac{T}{\lambda}\right) + \sum_{i=d+1}^{N} \ln \left(1 + \frac{t_i}{\lambda_{d+1}}\right) \\ &\leq d \ln \left(1 + \frac{T}{\lambda}\right) + \frac{T}{\lambda_{d+1}} \\ &\leq 2d \ln \left(1 + \frac{T}{\lambda}\right) \end{split}$$



#### SpectralUCB: Improving the running time

- Reduced basis: We only need first few eigenvectors.
- **Getting** J eigenvectors:  $\mathcal{O}(Jm \log m)$  time for m edges
- Computationally less expensive, comparable performance.



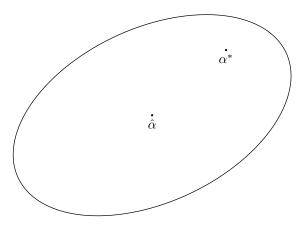


#### SpectralUCB: How to make it even faster?

- UCB-style algorithms need to (re)-compute UCBs every t
- Can be a problem for large set of arms  $\rightarrow D^2 N \rightarrow N^3$
- Optimistic (UCB) approach vs. Thompson Sampling
  - Play the arm maximizing probability of being the best
    - Sample  $\widetilde{\alpha}$  from the distribution  $\mathcal{N}(\widehat{\alpha}, v^2 \mathbf{V}^{-1})$
    - Play arm which maximizes  $\mathbf{x}^{\mathsf{T}} \widetilde{\boldsymbol{\alpha}}$  and observe reward
  - Compute posterior distribution according to reward received
- Only requires  $D^2 + DN \rightarrow N^2$  per step update

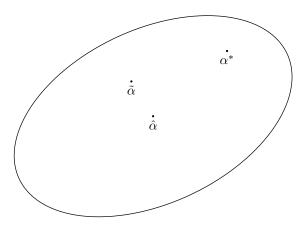


### **Thomson Sampling: Estimate**



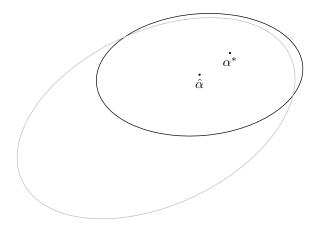


# **Thomson Sampling: Sample**



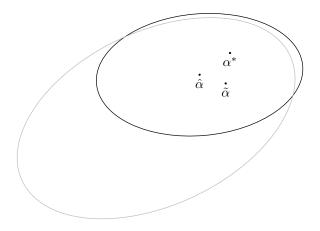


## **Thomson Sampling: Estimate**



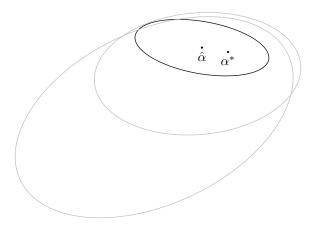


## **Thomson Sampling: Sample**





### Thomson Sampling: Estimate ....





#### **SpectralTS for Graphs**

1: Input: 2:  $N, T, \{\mathbf{A}_{\mathbf{L}}, \mathbf{Q}\}, \lambda, \delta, R, C$ 3: Initialization: 4:  $v = R\sqrt{6d \log((\lambda + T)/\delta\lambda)} + C$ 5:  $\widehat{\alpha} = 0_N$ 6:  $f = 0_N$ 7:  $\mathbf{V} = \mathbf{\Lambda}_{\mathbf{I}} + \lambda \mathbf{I}_{\mathbf{N}}$ 8: Run: 9: for t = 1 to T do 10: Sample  $\widetilde{\alpha} \sim \mathcal{N}(\widehat{\alpha}, v^2 \mathbf{V}^{-1})$ 11:  $\pi(t) \leftarrow \arg \max_{a} \mathbf{x}_{a}^{\mathsf{T}} \widetilde{\alpha}$ 12: Observe a noisy reward  $r(t) = \mathbf{x}_{\pi(t)}^{\mathsf{T}} \boldsymbol{\alpha}^* + \varepsilon_t$ 13:  $\mathbf{f} \leftarrow \mathbf{f} + \mathbf{x}_{\pi(t)} r(t)$ 14: Update  $\mathbf{V} \leftarrow \mathbf{V} + \mathbf{x}_{\pi(t)} \mathbf{x}_{\pi(t)}^{\mathsf{T}}$ 15: Update  $\widehat{\alpha} \leftarrow \mathbf{V}^{-1}\mathbf{f}$ 16: end for



#### SpectralTS: Regret bound

- ► *d*: Effective dimension.
- λ: Minimal eigenvalue of  $\Lambda = \Lambda_L + \lambda I$ .
- C: Smoothness upper bound,  $\|\alpha^*\|_{\mathbf{A}} \leq C$ .
- ►  $\mathbf{x}_i^{\mathsf{T}} \boldsymbol{\alpha}^* \in [-1, 1]$  for all *i*.

The **cumulative regret**  $R_T$  of **SpectralTS** is with probability  $1 - \delta$  bounded as

$$\mathcal{R}_{T} \leq \frac{11g}{p} \sqrt{\frac{4+4\lambda}{\lambda}} dT \log \frac{\lambda+T}{\lambda} + \frac{1}{T} + \frac{g}{p} \left(\frac{11}{\sqrt{\lambda}} + 2\right) \sqrt{2T \log \frac{2}{\delta}},$$
  
where  $p = 1/(4e\sqrt{\pi})$  and  
 $g = \sqrt{4\log TN} \left( R \sqrt{6d \log \left(\frac{\lambda+T}{\delta\lambda}\right)} + C \right) + R \sqrt{2d \log \left(\frac{(\lambda+T)T^{2}}{\delta\lambda}\right)} + C.$ 

$$R_T \approx d\sqrt{T \log N}$$



### SpectralTS: Analysis sketch

Divide arms into two groups

$$\Delta_i = \mathbf{x}_*^{\mathsf{T}} \boldsymbol{\alpha} - \mathbf{x}_i^{\mathsf{T}} \boldsymbol{\alpha} \leq g \|\mathbf{x}_i\|_{\mathbf{V}_*^{-1}}$$

arm *i* is **unsaturated** 

arm *i* is **saturated** 

#### Saturated arm

- Small standard deviation  $\rightarrow$  accurate regret estimate.
- ▶ High regret on playing the arm → Low probability of picking

#### Unsaturated arm

- Low regret bounded by a factor of standard deviation
- High probability of picking



#### SpectralTS: Analysis sketch

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- ▶ Confidence ellipsoid for estimate  $\widehat{\mu}$  of  $\mu$  (with probability  $1 \delta/T^2$ )
  - Using analysis of OFUL algorithm (Abbasi-Yadkori et al., 2011)

$$|\mathbf{x}_{i}^{\mathsf{T}}\widehat{\boldsymbol{\alpha}} - \mathbf{x}_{i}^{\mathsf{T}}\boldsymbol{\alpha}| \leq \left(R\sqrt{2\,d\log\left(\frac{(\lambda+T)\,T^{2}}{\delta\lambda}\right)} + C\right)\|\mathbf{x}_{i}\|_{\mathbf{V}_{t}^{-1}} = \ell\|\mathbf{x}_{i}\|_{\mathbf{V}_{t}^{-1}}$$

The key result coming from spectral properties of V<sub>t</sub>.

$$\log rac{|\mathbf{V}_t|}{|\mathbf{\Lambda}|} \leq 2d \log \left(1 + rac{T}{\lambda}
ight)$$

- ▶ Concentration of sample  $\widetilde{\alpha}$  around mean  $\widehat{\alpha}$  (with probability  $1 1/T^2$ )
  - Using concentration inequality for Gaussian random variable.

$$|\mathbf{x}_{i}^{\mathsf{T}}\widetilde{\boldsymbol{\alpha}} - \mathbf{x}_{i}^{\mathsf{T}}\widehat{\boldsymbol{\alpha}}| \leq \left(R\sqrt{6d\log\left(\frac{\lambda+T}{\delta\lambda}\right)} + C\right) \|\mathbf{x}_{i}\|_{\mathbf{V}_{t}^{-1}}\sqrt{4\log(TN)} = \nu \|\mathbf{x}_{i}\|_{\mathbf{V}_{t}^{-1}}\sqrt{4\log(TN)}$$



#### SpectralTS: Analysis sketch

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**Define** regret'(t) = regret(t)  $\cdot \mathbb{1}\{|\mathbf{x}_i^{\mathsf{T}}\widehat{\alpha}(t) - \mathbf{x}_i^{\mathsf{T}}\alpha| \le \ell \|\mathbf{x}_i\|_{\mathbf{V}_t^{-1}}\}$ 

$$\mathsf{regret}'(t) \leq \frac{11g}{p} \|\mathbf{x}_{\boldsymbol{a}(t)}\|_{\mathbf{V}_t^{-1}} + \frac{1}{\mathcal{T}^2}$$

Super-martingale (i.e.  $\mathbb{E}[Y_t - Y_{t-1}|\mathcal{F}_{t-1}] \leq 0$ )

$$\begin{aligned} X_t &= \operatorname{regret}'(t) - \frac{11g}{p} \|\mathbf{x}_{a(t)}\|_{\mathbf{V}_t^{-1}} - \frac{1}{T^2} \\ Y_t &= \sum_{w=1}^t X_w. \end{aligned}$$

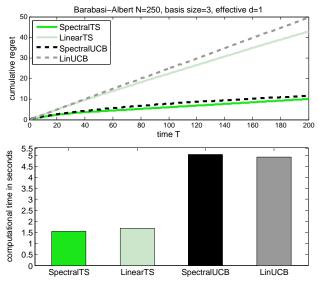
 $(Y_t; t = 0, ..., T)$  is a super-martingale process w.r.t. history  $\mathcal{F}_t$ .

Azuma-Hoeffding inequality for super-martingales, w.p.  $1-\delta/2$ :

$$\sum_{t=1}^{T} \operatorname{regret}'(t) \leq \frac{11g}{p} \sum_{t=1}^{T} \|\mathbf{x}_{a(t)}\|_{\mathbf{V}_{t}^{-1}} + \frac{1}{T} + \frac{g}{p} \left(\frac{11}{\sqrt{\lambda}} + 2\right) \sqrt{2T \ln \frac{2}{\delta}}$$



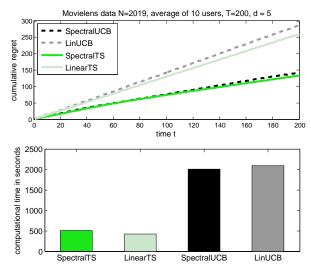
# Spectral Bandits: Synthetic experiment





#### Spectral Bandits: Real world experiment

MovieLens dataset of 6k users who rated one million movies.



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# **Spectral Bandits Summary**

- Spectral bandit setting (smooth graph functions).
- SpectralUCB

• Regret bound 
$$R_T = \widetilde{\mathcal{O}}\left(\frac{d}{\sqrt{T \ln T}}\right)$$

SpectralTS

• Regret bound 
$$R_T = \widetilde{O}\left(\frac{d}{\sqrt{T \ln N}}\right)$$

- Computationally more efficient.
- SpectralEliminator

• Regret bound 
$$R_T = \widetilde{O}\left(\sqrt{dT \ln T}\right)$$

- Better upper, empirically does not seem to work well (yet)
- Bounds scale with effective dimension  $d \ll D$ .



### SpectralEliminator: Pseudocode

Input:

N : the number of nodes, T : the number of pulls  $\{\Lambda_{I}, Q\}$  spectral basis of L  $\lambda$  : regularization parameter  $\beta$ ,  $\{t_i\}_i^J$  parameters of the elimination and phases  $A_1 \leftarrow \{\mathbf{x}_1, \ldots, \mathbf{x}_K\}.$ for i = 1 to J do  $\mathbf{V}_{t_i} \leftarrow \gamma \mathbf{\Lambda}_{\mathbf{L}} + \lambda \mathbf{I}$ for  $t = t_i$  to min $(t_{i+1} - 1, T)$  do Play  $\mathbf{x}_t \in A_i$  with the largest width to observe  $r_t$ :  $\mathbf{x}_t \leftarrow \arg \max_{\mathbf{x} \in A_i} \|\mathbf{x}\|_{\mathbf{V}^{-1}}$  $\mathbf{V}_{t+1} \leftarrow \mathbf{V}_t + \mathbf{x}_t \mathbf{x}_t^{\mathsf{T}}$ end for Eliminate the arms that are not promising:  $\widehat{\alpha}_t \leftarrow \mathbf{V}_t^{-1}[\mathbf{x}_{t_i}, \dots, \mathbf{x}_t][r_{t_i}, \dots, r_t]^{\mathsf{T}}$  $A_{j+1} \leftarrow \left\{ \mathbf{x} \in A_j, \langle \widehat{\boldsymbol{\alpha}}_t, \mathbf{x} \rangle + \|\mathbf{x}\|_{V_{\bullet}^{-1}} \beta \ge \max_{\mathbf{x} \in A_j} \left[ \langle \widehat{\boldsymbol{\alpha}}_t, \mathbf{x} \rangle - \|\mathbf{x}\|_{V_{\bullet}^{-1}} \beta \right] \right\}$ end for

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#### SpectralEliminator: Analysis

#### SpectralEliminator

 Divide time into sets (t<sub>1</sub> = 1 ≤ t<sub>2</sub> ≤ ...) to introduce independence for Azuma-Hoeffding inequality and observe
 R<sub>T</sub> ≤ ∑<sub>j=0</sub><sup>J</sup>(t<sub>j+1</sub> − t<sub>j</sub>)[⟨**x**<sup>\*</sup> − **x**<sub>t</sub>, â<sub>j</sub>⟩ + (||**x**<sup>\*</sup>||<sub>V<sub>j</sub><sup>-1</sup></sub> + ||**x**<sub>t</sub>||<sub>V<sub>j</sub><sup>-1</sup></sub>)β]

• Bound 
$$\langle \mathbf{x}^* - \mathbf{x}_t, \widehat{oldsymbol{lpha}}_j 
angle$$
 for each phase

- $\blacktriangleright \text{ No bad arms: } \langle \mathbf{x}^* \mathbf{x}_t, \widehat{\alpha}_j \rangle \leq (\|\mathbf{x}^*\|_{\mathbf{V}_j^{-1}} + \|\mathbf{x}_t\|_{\mathbf{V}_j^{-1}})\beta$
- By algorithm:  $\|\mathbf{x}\|_{\mathbf{V}_{j}^{-1}}^{2} \leq \frac{1}{t_{j}-t_{j-1}} \sum_{s=t_{j-1}+1}^{t_{j}} \|\mathbf{x}_{s}\|_{\mathbf{V}_{s-1}^{-1}}^{2}$

$$\blacktriangleright \sum_{s=t_{j-1}+1}^{t_j} \min\left(1, \|\mathbf{x}_s\|_{\mathbf{V}_{s-1}}^2\right) \le \log \frac{|\mathbf{V}_j|}{|\mathbf{A}|}$$



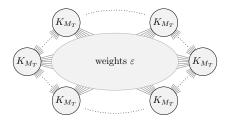
# Spectral Bandits: Is it possible to do better?

Is d a good quantity that embodies the difficulty?

Lower bound!

For any d, we construct a graph that for any reasonable algorithm, the regret is at least  $\Omega(\sqrt{dT})$ .

How? By reduction to *d*-arm bandits problem.





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