# inventors for the digital world

### THE POWER OF GRAPHS IN SPEEDING UP ONLINE LEARNING AND DECISION MAKING



#### Michal Valko

SequeL @ Inria Lille - Nord Europe







Berkeley's floating sensor network



Erdös number project









### MY PAST 10 YEARS WITH GRAPH IN ML



Online semi-supervised learning for personalization

Bandits and MDPs with discrete and continuous variables

### JOINT WORK WITH...





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**Gergely Neu** 

U Pompeu Fabra

Zheng Wen Adobe Research



**Eugene Belilovsky** MILA 7 Sequen





# ONLINE LEARNING

when we reason on the fly



### IN 2007 IT ALL STARTED WITH AN IDEA...

- Develop sequential machine learning recognition system
- System with **minimal feedback**
- 90% accurate over 90% of time
- With **theory** that guarantee's its performance
- Efficient (e.g., mobile device)









from B. Kveton

### ... AND RESULTED IN A REAL SYSTEM IN 2009

- adaptive graph-based recognition system
  - highly accurate
  - trained from a small amount of labeled data
  - real-time running time
  - robust to outliers
  - theoretical analysis

$$\frac{1}{n} \sum_{t} \left( \ell_t^{\mathsf{q}}[t] - y_t \right)^2 \le \frac{1}{n_l} \sum_{i \in I} \left( l_i^* - y_i \right)^2 + \mathcal{O}(n^{-\frac{1}{2}})$$





from B. Kveton





<u>Conditional</u> anomalies are often medical errors. "Medical errors account for 200 000 preventable deaths a year." (HealthGrades study, Wall Street Journal, July 27<sup>th</sup> 2004)





Online Semi-Supervised Learning on Quantized Graphs. In Proceedings of the 26th Conference on Uncertainty in Artificial Intelligence, California, July 2010.

### FACE-RECOGNITION FOR INTEL



### THIS CAN'T SCALE: CONNECTED CAR





#### Personalization





### SIZE and SPEED

### ANOMALIES





15 Sequel

### HUGE AND/OR ONLINE





MV, Kveton, Huang, Ting: Online Semi-Supervised Learning on Quantized Graphs UAI 2010
Kveton, MV, Rahimi, Huang: Semi-Supervised Learning with Max-Margin Graph Cuts AISTATS 2010
Calandriello, Lazaric, MV: Distributed sequential sampling for kernel matrix approximation AISTATS 2017
Calandriello, Lazaric, MV: Second-order kernel online convex optimization with adaptive sketching, ICML 2017
Calandriello, Lazaric, MV: Efficient second-order online kernel learning with adaptive embedding, NIPS 2017
Calandriello, Koutis, Lazaric, MV: Improved large-scale graph learning through ridge spectral sparsification, ICML 2018
code: <a href="http://researchers.lille.inria.fr/~valko/hp/publications/squeak.py">http://researchers.lille.inria.fr/~valko/hp/publications/squeak.py</a>

## Industry transfer to (intel)

#### Context Aware Vehicle

recognizes when your face is turned to the side

#### Everyday Sensing and Perception

- health monitoring and assisted living
- Google TV project
  - Personalized advertisement
- Connected Cars
  - Ford, Toyota, Audi/VW Group, Nissan
- Intel Phone (marketed in 2015)
  - adaptive logging in









6870 lines of code in C++ using OpenCV library 2-3 years of research + development

## Technology transfer to UPMC (2011)

3 NIH grants \$2,961,032

14 GB of data, 27667 lines of code, 2007-2011.

Homer Warner Award 2010

#### Example: Heparin Induced Thrombocytopenia

- BEFORE: about 10 years of creating the rule
- BEFORE: Rule definition has 5 pages
- BEFORE: Every adjustment takes 3 months
- AFTER: 5 years of historical data (no supervision needed)
- AFTER: Better performance (prediction/recall) than for the rule
- Large study: 734 decisions (orders) for 40 000 cases
- Evaluation: 54.5% of alerts found useful
- Used by Department of Clinical Care
- Explainability

# ONLINE DECISION-MAKING

when we want to act





### Example of a graph bandit problem

#### movie recommendation

- recommend movies to a single user
- goal: maximise the sum of the ratings (minimise regret)
- good prediction after just a few steps
  - $T \ll N$
- extra information
  - ratings are smooth on a graph
- main question: can we learn faster?

### **GETTING REAL**



#### Let's be lazy and ignore the structure



Multi-armed bandit problem!

Worst case regret (to the best fixed strategy)

Matching lower bound (Auer, Cesa-Bianchi, Freund, Schapire 2002)

**How big is N?** Number of movies on <u>http://www.imdb.com/stats</u>: 5,310,913 **Problem:** Too many actions!





 $R_T = \mathcal{O}\left(\sqrt{NT}\right)^{\#r}$ 

- Arm independence is too strong and unnecessary
- Replace N with something much smaller

**LEARNING FASTER** 

- problem/instance/data dependent
- example: linear design N to D
- Here use Graphs to encode structure of decision making!
  - sequential problems where actions are nodes on a graph
  - ▶ find strategies that replace **N** with a **smaller graph-dependent** quantity







**#rounds** 

#dimensions



### **GRAPH BANDITS: GENERAL SETUP**

#### Every round **t** the learner

- ▶ picks a node  $I_t \in [N]$
- ▶ incurs a loss  $\ell_{t,I_t}$
- optional feedback

The performance is total expected regret

$$R_{T} = \max_{i \in [N]} \mathbb{E} \left[ \sum_{t=1}^{T} (\ell_{t,I_{t}} - \ell_{t,i}) \right]$$

### STRUCTURES IN ONLINE (RL/BANDIT) PROBLEMS

### GRAPHS

### **KERNELS**

### **DISCOUNT FACTOR in MDPs**

### **CONTINUOUS FUNCTIONS**

### **STRUCTURES WITHOUT TOPOLOGY**









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### **SPECIFIC GRAPH BANDIT SETTINGS**





MV, Munos, Kveton, Kocák: **Spectral Bandits for Smooth Graph Functions**, ICML 2014 Kocák, MV, Munos, Agrawal: **Spectral Thompson Sampling**, AAAI 2014 Hanawal, Saligrama, MV, Munos: **Cheap Bandits**, ICML 2015

# SPECTRAL BANDITS

exploiting smoothness of rewards on graphs



### **SPECTRAL BANDITS**

#### Assumptions

- Unknown reward function  $f : V(G) \rightarrow \mathbb{R}$ .
- Function f is smooth on a graph.
- Neighboring movies  $\Rightarrow$  similar preferences.
- Similar preferences  $\neq$  neighboring movies.

### **C**Desiterata

An algorithm useful in the case  $T \ll N!$ 



### **FLIXSTER DATA**





Eigenvectors from the Flixster data corresponding to the smallest few eigenvalues of the graph Laplacian projected onto the first principal component of data. Colors indicate the values.

#### 29 Segues

### LINEAR VS. SPECTRAL BANDITS

- Linear bandit algorithms
  - ► LinUCB
    - Regret bound  $\approx D\sqrt{T \ln T}$
  - ► LinearTS
    - Regret bound  $\approx D\sqrt{T \ln N}$

(Agrawal and Goyal, 2013)

(Li et al., 2010)

**Note:** *D* is ambient dimension, in our case *N*, length of  $x_i$ . Number of actions, e.g., all possible movies  $\rightarrow$  **HUGE!** 

Spectral bandit algorithms

- SpectralUCB
  - Regret bound  $\approx d\sqrt{T \ln T}$
  - Operations per step:  $D^2N$
- SpectralTS

(Kocák et al., AAAI 2014)

(Valko et al., ICML 2014)

- Regret bound  $\approx d\sqrt{T \ln N}$
- Operations per step:  $D^2 + DN$

**Note:** *d* is effective dimension, usually much smaller than *D*.





### **SPECTRALUCB REGRET BOUND**

- ► *d*: Effective dimension.
- >  $\lambda$ : Minimal eigenvalue of  $\Lambda = \Lambda_L + \lambda I$ .

cumulative <sup>20</sup>

ETTER

- C: Smoothness upper bound,  $\| \alpha^* \|_{\Lambda} \leq C$ .
- ►  $\mathbf{x}_i^{\mathsf{T}} \boldsymbol{\alpha}^* \in [-1, 1]$  for all *i*.

The **cumulative regret**  $R_T$  of **SpectralUCB** is with probability  $1 - \delta$  bounded as

$$R_{T} \leq \left(8R\sqrt{d\ln\frac{\lambda+T}{\lambda}} + 2\ln\frac{1}{\delta} + 4C + 4\right)\sqrt{dT\ln\frac{\lambda+T}{\lambda}}.$$





Kocák, Neu, MV, Munos: **Efficient learning by implicit exploration in bandit problems with side observations**, NIPS 2014

Kocák, Neu, MV: **Online learning with Erdos-Rényi side-observation graphs** UAI 2016

Kocák, Neu, MV: Online learning with noisy side observations, AISTATS 2016

GRAPH BANDITS WITH SIDE **OBSERVATIONS** exploiting free observations from neighbouring nodes



### **SIDE OBSERVATIONS: UNDIRECTED**



#### **Example 1: undirected observations**



### **SIDE OBSERVATIONS: DIRECTED**



#### **Example 2: Directed observation**





#### **Full-information**

- observe losses of all actions
- example: Hedge

 $R_T = \widetilde{\mathcal{O}}(\sqrt{T})$ 

#### **Bandits**

- observe losses of the chosen action
- example: EXP3

$$\mathsf{R}_{\mathcal{T}} = \widetilde{\mathcal{O}}(\sqrt{\mathsf{NT}})$$



### **KNOWLEDGE OF OBSERVATION GRAPHS**

- ELP (Mannor and Shamir 2011)
  - EXP3 with "LP balanced exploration"
  - undirected  $O(\sqrt{\alpha T})$   $\Box$ -needs to know  $G_t$
  - directed case  $O(\sqrt{(cT)})$  needs to know  $G_t$
- EXP3-SET (Alon, Cesa-Bianchi, Gentile, Mansour, 2013)
  - undirected  $O(\mathcal{J}(\alpha T))$   $\square$  does not need to know  $G_t$   $\square$
- EXP3-DOM (Alon, Cesa-Bianci Innia le, mansour, 2013)
  - directed  $O(\sqrt{\alpha T})$   $\nabla$ -need to know  $G_t$
  - calculates dominating set







### **EXP3-IX: IMPLICIT EXPLORATION**

Algorithm 1 EXP3-IX Benefits of the **implicit exploration** 1: Input: Set of actions S = [d], parameters  $\gamma_t \in (0, 1), \eta_t > 0$  for  $t \in [T]$ . 2: no need to know the graph before 3: for t = 1 to T do  $w_{t,i} \leftarrow (1/d) \exp\left(-\eta_t \widehat{L}_{t-1,i}\right)$  for  $i \in [d]$ 4: An adversary privately chooses losses  $\ell_{t,i}$ 5: no need to estimate dominating set for  $i \in [d]$  and generates a graph  $G_t$  $W_t \leftarrow \sum_{i=1}^d w_{t,i}$ 6: no need for doubling trick  $p_{t,i} \leftarrow w_{t,i}/W_t$ 7: Choose  $I_t \sim p_t = (p_{t,1}, ..., p_{t,d})$ 8: Observe graph  $G_t$ 9: no need for aggregation Observe pairs  $\{i, \ell_{t,i}\}$  for  $(I_t \to i) \in G_t$ 10:  $o_{t,i} \leftarrow \sum_{(j \to i) \in G_t} p_{t,j} \text{ for } i \in [d]$ 11:  $\hat{\ell}_{t,i} \leftarrow \frac{\ell_{t,i}}{o_{t,i} + \gamma_t} \mathbb{1}_{\{(I_t \to i) \in G_t\}} \text{ for } i \in [d]$ 12:  $R_T = \widetilde{\mathcal{O}}\left(\sqrt{\overline{\alpha}\,T\,\ln\,N}\right)$ 13: **end for** ······

Optimistic bias for the loss estimates

$$\mathbb{E}[\hat{\ell}_{t,i}] = \frac{\ell_{t,i}}{o_{t,i} + \gamma} o_{t,i} + 0(1 - o_{t,i}) = \ell_{t,i} - \ell_{t,i} \frac{\gamma}{o_{t,i} + \gamma} \leq \ell_{t,i}$$



### FOLLOW UPS

- EXP3-IX (Kocák, Neu, MV, Munos, 2014)
  - directed O(√(αT)) ✓ does not need to know Gt ✓
- EXP3.G (Alon, Cesa-Bianchi, Dekel, Koren, 2015)
  - directed  $O(\int(\alpha T)) \Box$  does not need to know  $G_t \Box$
  - mixes uniform distribution
  - more general algorithm for settings beyond bandits

Inría

- high-probability bound *(mia-*)
- Neu 2015: high-probability bound for EXP3-IX
- ▶ TextBook: Bandit Algorithms T. Lattimore & Cs. Szepesvári







### **EXTENSION: COMPLEX GRAPH ACTIONS**



**Example:** online shortest path semi-bandits with observing traffic on the side streets



- ▶ Play action  $\mathbf{V}_t \in S \subset \{0,1\}^N$ ,  $\|\mathbf{v}\|_1 \leq m$  from all  $\mathbf{v} \in S$
- ► Obtain losses  $\mathbf{V}_t^{\mathsf{T}} \boldsymbol{\ell}_t$
- Observe additional losses according to the graph

$$R_{T} = \widetilde{\mathcal{O}}\left(m^{3/2}\sqrt{\sum_{t=1}^{T}\alpha_{t}}\right) = \widetilde{\mathcal{O}}\left(m^{3/2}\sqrt{\overline{\alpha}T}\right)$$



### **EXTENSION: NOISY SIDE OBSERVATIONS**





Want: only reliable information!

1) If we know the perfect cutoff  $\epsilon$ 

- reliable: use as exact
- unreliable: rubbish

then we can improve over pure bandit setting!2) Treating noisy observation induces bias

What can we hope for?

$$\widetilde{\mathcal{O}}\left(\sqrt{\mathbf{1}T}\right) \leq \widetilde{\mathcal{O}}\left(\sqrt{\overline{\alpha}^{\star}T}\right) \leq \widetilde{\mathcal{O}}\left(\sqrt{\mathbf{N}T}\right)$$

effective independence number

Can we learn without knowing either  $\epsilon$  or  $\alpha^*$ ?



### **NEW DIRECTIONS: UNKNOWN GRAPHS!**

- Learning on the graph while learning the graph?
  - most of algorithms require (some) knowledge of the graph
  - not always available to the learner
- Question: Can we learn faster without knowing the graphs?
  - example: social network provider has little incentive to reveal the graphs to advertisers
- Answer: Cohen, Hazan, and Koren: Online learning with feedback graphs without the graphs (ICML 2016)
  - **NO!** (in general we cannot, but possible in the stochastic case)
- ▶ NEXT in the talk: examples where we can do something!
  - Erdös-Rényi side observation graphs
  - Influence Maximisation

### **EXTENSION: ERDÖS-RÉNYI GRAPHS**





- N-2 samples from Bernoulli(r<sub>t</sub>) ... R(k)
- N-2 samples from pti ... P(k)
- ▷ O'(k) = P(k) + (1-P(k))R(k)
- $G_{ti} = \min\{k: O'(k) = 1\} \cup \{N-1\}$

$$E[G_{ti}] \approx 1/(p_{ti} + (1-p_{ti})r_t)$$
$$\widehat{\ell}_{t,i} = G_{t,i}O_{t,i}\ell_{t,i}$$

If  $r_t \ge (\log T)/(2N-2)$  then  $\mathcal{O}\left(\sqrt{\log N \sum_{t=1}^T \frac{1}{r_t}}\right)$ 

Lower bound (Alon et al. 2013)  $\Omega(\sqrt{T/r})$ Get rid of rt  $\geq (\log T)/(2N-2)$ ? Carpentier, MV: **Revealing Graph Bandits for Maximising Local Influence**, AISTATS 2016 Wen, Kveton, MV: **Influence Maximization with Semi-Bandit Feedback**, NIPS 2017

# INFLUENCE MAXIMISATION looking for the influential nodes while exploring the graph



### **REVEALING BANDITS FOR LOCAL INFLUENCE**

Unknown  $\mathbf{M} = (p_{i,j})_{i,j}$  symmetric matrix of influences

In each time step  $t = 1, \ldots, T$ 

- $\blacktriangleright$  learners picks a node  $k_t$
- ▶ set  $S_{k_t,t}$  of influenced nodes is *revealed*

Ínría Select influential people = Find the strategy maximising  $L_T = \sum |S_{k_t,t}|$ 

The number of expected influences of node **k** is by definition

$$r_k = \mathbb{E}\left[|S_{k,t}|\right] = \sum_{j \le N} p_{k,j}$$

Oracle strategy always selects the best

$$k^* = \arg\max_k \mathbb{E}\left[\sum_{t=1}^T |S_{k,t}|\right] = \arg\max_k Tr_k$$

Expected regret of any adaptive, non-oracle strategy unaware of M

 $\mathbb{E}[R_T] = \mathbb{E}[L_T^*] - \mathbb{E}[L_T]$ 





## GLOBAL DIFFUSION PROCESS OF A MARKOV CHAIN



- Sets of progressive diffusion
  - modeling diffusion steps
- Random stopping time
  - but bounded
- Topological ordering

$$\mathcal{S}^{0} \stackrel{\Delta}{=} \mathcal{S}_{t}$$
$$\mathcal{S}^{\tau+1} \stackrel{\Delta}{=} \left\{ u_{2} \in \mathcal{V}_{\mathcal{S}_{t},v} : u_{2} \notin \bigcup_{\tau'=0}^{\tau} \mathcal{S}^{\tau'} \text{ and } \exists e = (u_{1}, u_{2}) \in \mathcal{E}_{\mathcal{S}_{t},v} \text{ s.t. } u_{1} \in \mathcal{S}^{\tau} \text{ and } \mathbf{w}(e) = 1 \right\}$$

### **EMPIRICAL RESULTS**





### **GLOBAL WORST-CASE BOUNDS**





| topology       | $C_{\mathcal{G}}$ (worst-case $C_*$ )  | $R^{lpha\gamma}(n)$ for general ${f X}$  | $R^{\alpha\gamma}(n)$ for $\mathbf{X} = \mathbf{I}$                           |
|----------------|--|--|---|
| bar graph      | $\mathcal{O}(\sqrt{K})$                | $\widetilde{\mathcal{O}}\left( dK\sqrt{n}/(lpha\gamma) ight)$                  | $\widetilde{\mathcal{O}}\left(L\sqrt{Kn}/(lpha\gamma) ight)$                  |
| star graph     | $\mathcal{O}(L\sqrt{K})$               | $\widetilde{\mathcal{O}}\left(dL^{\frac{3}{2}}\sqrt{Kn}/(lpha\gamma) ight)$    | $\widetilde{\mathcal{O}}\left(L^2\sqrt{Kn}/(lpha\gamma) ight)$                |
| ray graph      | $\mathcal{O}(L^{\frac{5}{4}}\sqrt{K})$ | $\widetilde{\mathcal{O}}\left(dL^{\frac{7}{4}}\sqrt{Kn}/(\alpha\gamma)\right)$ | $\widetilde{\mathcal{O}}\left(L^{\frac{9}{4}}\sqrt{Kn}/(\alpha\gamma)\right)$ |
| tree graph     | $\mathcal{O}(L^{\frac{3}{2}})$         | $\widetilde{\mathcal{O}}\left(dL^2\sqrt{n}/(lpha\gamma) ight)$                 | $\widetilde{\mathcal{O}}\left(L^{\frac{5}{2}}\sqrt{n}/(\alpha\gamma)\right)$  |
| grid graph     | $\mathcal{O}(L^{\frac{3}{2}})$         | $\widetilde{\mathcal{O}}\left(dL^2\sqrt{n}/(lpha\gamma) ight)$                 | $\widetilde{\mathcal{O}}\left(L^{\frac{5}{2}}\sqrt{n}/(\alpha\gamma)\right)$  |
| complete graph | $\mathcal{O}(L^2)$                     | $\widetilde{\mathcal{O}}\left(dL^3\sqrt{n}/(lpha\gamma) ight)$                 | $\widetilde{\mathcal{O}}\left(L^4\sqrt{n}/(lpha\gamma) ight)$                 |

Table 1:  $C_{\mathcal{G}}$  and *worst-case* regret bounds for different graph topologies





- real Facebook (a small subgraph)
- weights from U(0,0.1)
- nodetovec with d=10
  - imperfect
- ▶ K = 10
- CUCB with no linear generalisation





- Active learning on graphs: online influence maximization
  - learning the graph while acting on it optimally
  - global cascading model with edge level feedback
  - **difficulty of the problem** and scaling with it
- What is next?
  - node-level feedback
  - dynamic/evolving graphs
  - realistic accessibility constraints



#### Survey: <a href="http://researchers.lille.inria.fr/~valko/hp/publications/valko2016bandits.pdf">http://researchers.lille.inria.fr/~valko/hp/publications/valko2016bandits.pdf</a>

48 Segues

### SCALE UP!!!





MV, Kveton, Huang, Ting: **Online Semi-Supervised Learning on Quantized Graphs** UAI 2010 Kveton, MV, Rahimi, Huang: **Semi-Supervised Learning with Max-Margin Graph Cuts** AISTATS 2010 Calandriello, Lazaric, MV: **Distributed sequential sampling for kernel matrix approximation** AISTATS 2017 Calandriello, Lazaric, MV: **Second-order kernel online convex optimization with adaptive sketching**, ICML 2017 Calandriello, Lazaric, MV: **Efficient second-order online kernel learning with adaptive embedding**, NIPS 2017 Calandriello, Koutis, Lazaric, MV: **Improved large-scale graph learning through ridge spectral sparsification**, ICML 2018 **code:** <u>http://researchers.lille.inria.fr/~valko/hp/publications/squeak.py</u>



- Large graphs do not fit in a single machine memory
- multiple passes slow, distribution has communication costs
- removing edges impacts structure/accuracy
- Make the graph sparse, while preserving its structure for learning

$$(1-\varepsilon)\mathsf{L}_{\mathcal{G}} \preceq \mathsf{L}_{\mathcal{H}} \preceq (1+\varepsilon)\mathsf{L}_{\mathcal{G}}$$

$$(1-\varepsilon)\mathsf{L}_{\mathcal{G}} - \varepsilon\gamma\mathsf{I} \preceq \mathsf{L}_{\mathcal{H}} \preceq (1+\varepsilon)\mathsf{L}_{\mathcal{G}} + \varepsilon\gamma\mathsf{I}$$

Mixed multiplicative / additive error

large (i.e.  $\geq \gamma$ ) directions reconstructed accurately small (i.e.  $\leq \gamma$ ) directions uniformly approximated ( $\gamma \mathbf{I}$ )

### **HOW DOES IT WORK?**





arbitrarily split in subgraphs that fit in a single machine

Ínría

### **DISRE GUARANTEES**





#### Theorem

Given an arbitrary graph  $\mathcal{G}$  w.h.p. DISRE satisfies

(1) each sub-graphs is an  $(\varepsilon, \gamma)$ -sparsifier

(2) with at most  $\mathcal{O}(d_{\text{eff}}(\gamma)\log(n))$  edges.



Dataset: Amazon co-purchase graph [Yang and Leskovec 2015]
↓ natural, artificially sparse (true graph known only to Amazon)
↓ we compute 4-step random walk to recover removed co-purchases
[Gleich and Mahoney 2015]

**Target:** eigenvector **v** associated with  $\lambda_2(\mathbf{L}_{\mathcal{G}})$  [Sadhanala et al. 2016]

n = 334,863 nodes, m = 98,465,352 edges (294 avg. degree)

| Alg.  | Parameters        | $ \mathcal{E} $ (x10 <sup>6</sup> ) | $\ \widetilde{\mathbf{f}}-\mathbf{v}\ _2^2 \ (\sigma\!=\!10^{-3})$ | $\ \widetilde{\mathbf{f}}-\mathbf{v}\ _2^2 (\sigma = 10^{-2})$ |
|-------|-------------------|-------------------------------------|--|--|
| EXACT |                   | 98.5                                | $0.067 \pm 0.0004$   | $0.756\pm0.006$  |
| kN    | k = 60            | 15.7                                | $0.172\pm0.0004$   | $0.822\pm0.002$  |
| DISRE | $\gamma\!=\!$ 0   | 22.8                                | $0.068\pm0.0004$   | $0.756 \pm 0.005$  |
| DISRE | $\gamma\!=\!10^2$ | 11.8                                | $\textbf{0.068} \pm 0.0002$  | $0.772 \pm 0.004$  |

**Time:** Loading  $\mathcal{G}$  from disk 90sec, DISRE 120sec( $k = 4 \times 32$  CPU), computing  $\tilde{\mathbf{f}}$  120sec, computing  $\hat{\mathbf{f}}$  720sec

### **CONCLUSION AND NEXT STEPS**



- Graphs give way to reason and act with relationships
  - step above over only considering entities
  - high-level cognition this is how children learn!
  - online learning and online decision-making
  - optimal allocation of resources (samples, time)
  - trees and Monte-Carlo tree search
  - tools to scale up the learning with near-linear time!
- What is next?
  - find the way to low-level representation
  - graph-networks that operate on the graphs and relational networks
  - intrinsic exploration over graph and other structures





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