

Quotients of the magmatic operad:
lattice structures and convergent rewrite systems

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I. Motivations

- ▷ Motivating example: an oscillating Hilbert series
- ▷ Nonsymmetric operads
- ▷ Presentations and Gröbner bases for operads

II. Magmatic quotients

- ▷ The category of magmatic quotients
- ▷ The lattice of magmatic quotients
- ▷ A Grassmann formula analog

III. Comb associative operads

- ▷ Definition of CAs operads
- ▷ The lattice of CAs operads
- ▷ Completion of CAs operads

IV. Conclusion and perspectives

Plan

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- ▶ Is the operad $\mathbf{CAs}^{(3)}$ presented by a finite Gröbner basis?
 - ▷ Yes: using the Buchberger/Knuth-Bendix's completion procedure.

- A **nonsymmetric linear operad** is a positively graded (\mathbb{K} -)vector space

$$\mathcal{O} = \bigoplus_{n \in \mathbb{N}} \mathcal{O}(n),$$

together with

- a distinguished element $\mathbf{1} \in \mathcal{O}(1)$;
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- ▶ How to construct operads?

- ▶ Using **presentations** by **generators and relations** $\langle \mathcal{X} \mid \mathcal{R} \rangle$.

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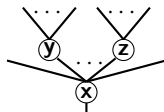
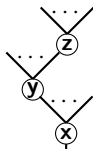


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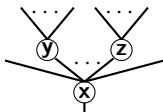
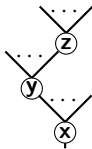


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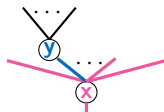


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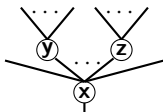
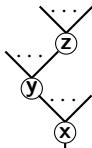
▷ $x \circ_i y$: obtained by grafting the root of y on the i -th leaf of x .

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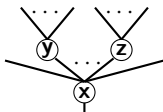
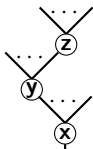
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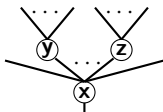
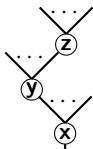
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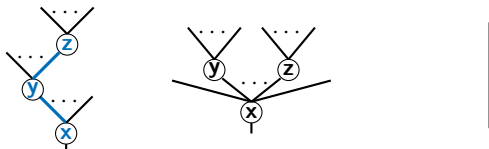
▷ neutrality of $\mathbf{1}$ for each \circ_i : $\mathbf{1} \circ_1 x = x = x \circ_i \mathbf{1}$;

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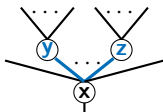
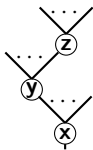
▷ **associativity** of **sequential compositions**: $x \circ_i (y \circ_j z) = (x \circ_i y) \circ_{i+j-1} z$;

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► **commutativity** of **parallel compositions**: $(x \circ_i y) \circ_{j+m-1} z = (x \circ_j z) \circ_i y$, where $i < j$ and m is the arity of y .

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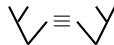


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 - ▷ one 0-ary generator (\rightsquigarrow the unit) and one binary generator (\rightsquigarrow the multiplication);



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- ▶ 2nd example: the **differential associative operad** is presented by

- ▶ Given $\mathcal{R} \subseteq \mathcal{F}(\mathcal{X})$, the operad presented by $\langle \mathcal{X} \mid \mathcal{R} \rangle$ is constructed as follows:
 - ▷ $\equiv_{\mathcal{R}}$: the operadic congruence generated by \mathcal{R} , that is $\mathbf{x} \equiv_{\mathcal{R}} 0$ for every $\mathbf{x} \in \mathcal{R}$;
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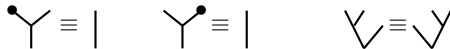


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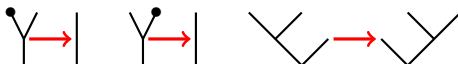
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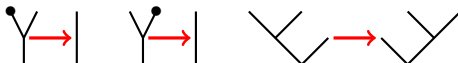
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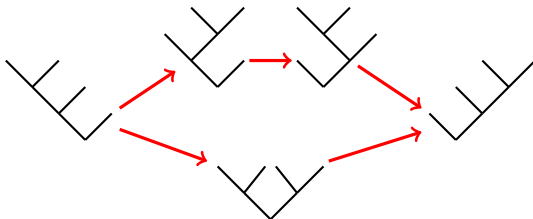
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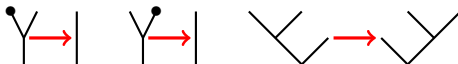
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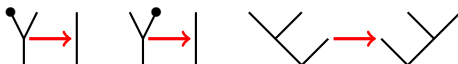
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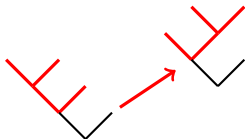
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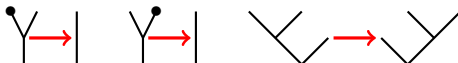
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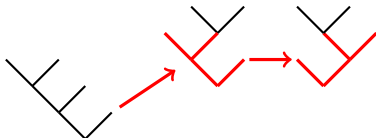
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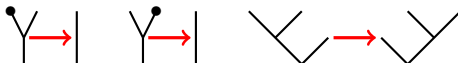
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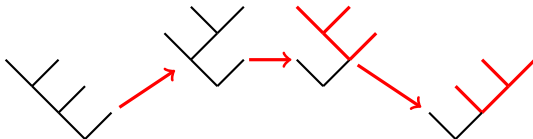
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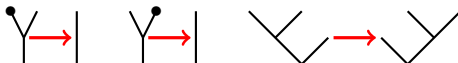
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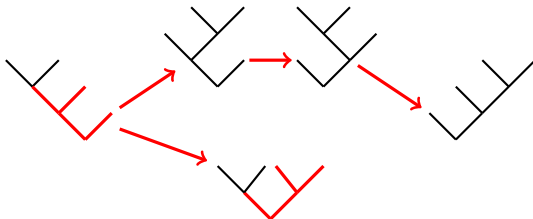
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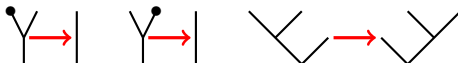
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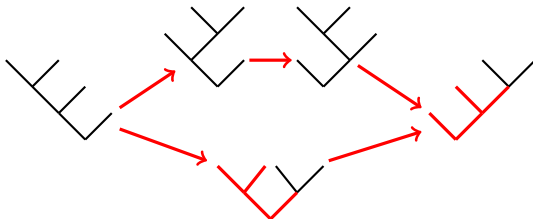
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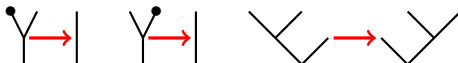
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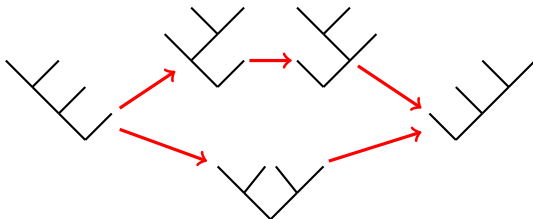
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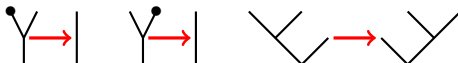
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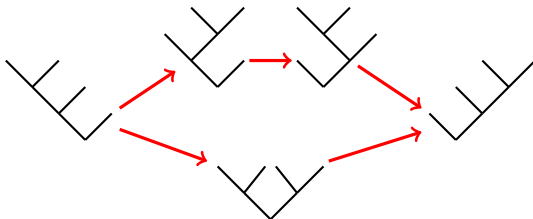
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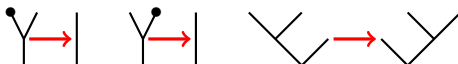


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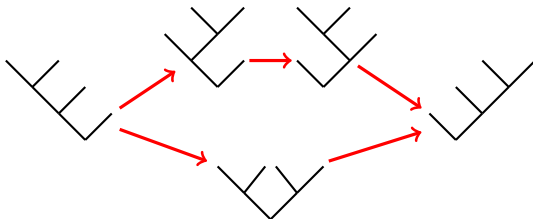


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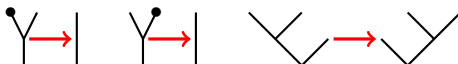
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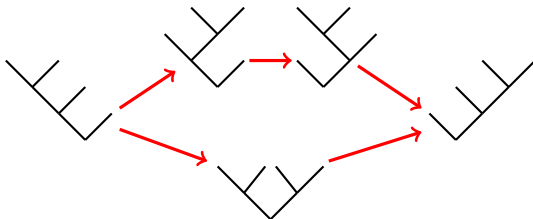
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 - ▶ the operad $\mathbf{CAs}^{(3)}$ belongs to a set of operads $\mathbf{CAs} := \{\mathbf{CAs}^{(\gamma)} \mid \gamma \geq 1\}$;
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- ▶ We study the induced poset on \mathbf{CAs} :
 - ▶ we present **new** lattice operations on this poset;
 - ▶ we study the existence of finite Gröbner bases for $\mathbf{CAs}^{(\gamma)}$ operads;
 - ▶ we deduce the complete expression of the Hilbert series of $\mathbf{CAs}^{(3)}$.

Plan

II. Lattice of magmatic quotients

- ▶ \mathbb{K} : a fixed field s.t. $\text{char}(\mathbb{K}) \neq 2$.
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Remark. A nonzero operad morphism between magmatic quotients is surjective.

► Let $\mathcal{O}_1 = \mathbb{K}\mathbf{Mag}/l_1$ and $\mathcal{O}_2 = \mathbb{K}\mathbf{Mag}/l_2$;

- ▷ we have $\dim(\mathrm{Hom}(\mathcal{O}_1, \mathcal{O}_2)) \leq 1$;
- ▷ $\dim(\mathrm{Hom}(\mathcal{O}_1, \mathcal{O}_2)) = 1$ iff $l_1 \subseteq l_2$;
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- ▶ Let $\preceq_i \subsetneq \mathcal{Q}(\mathbb{K}\mathbf{Mag}) \times \mathcal{Q}(\mathbb{K}\mathbf{Mag})$ defined by
 - ▷ $\mathcal{O}_2 \preceq_i \mathcal{O}_1$ iff $\dim(\mathrm{Hom}(\mathcal{O}_1, \mathcal{O}_2)) = 1$;

- ▶ Let $\theta_1 = \mathbb{K}\mathbf{Mag}/_{h_1}$ and $\theta_2 = \mathbb{K}\mathbf{Mag}/_{h_2}$;
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Theorem [C.-Cordero-Giraud, 2018]. Consider the notations introduced above.

- i. The tuple $(\mathcal{Q}(\mathbb{K}\mathbf{Mag}), \preceq_i, \wedge_i, \vee_i)$ is a lattice.

- ▶ Let $\theta_1 = \mathbb{K}\mathbf{Mag}/_{h_1}$ and $\theta_2 = \mathbb{K}\mathbf{Mag}/_{h_2}$;
 - ▷ we have $\dim(\mathrm{Hom}(\theta_1, \theta_2)) \leq 1$;
 - ▷ $\dim(\mathrm{Hom}(\theta_1, \theta_2)) = 1$ iff $h_1 \subseteq h_2$;
 - ▷ $\dim(\mathrm{Hom}(\theta_1, \theta_2)) = 1$ iff $\exists \varphi : \theta_1 \rightarrow \theta_2$ surjective.

- ▶ Let $\preceq_i \subseteq \mathcal{Q}(\mathbb{K}\mathbf{Mag}) \times \mathcal{Q}(\mathbb{K}\mathbf{Mag})$ defined by
 - ▷ $\theta_2 \preceq_i \theta_1$ iff $\dim(\mathrm{Hom}(\theta_1, \theta_2)) = 1$;

- ▶ Let $\wedge_i, \vee_i : \mathcal{Q}(\mathbb{K}\mathbf{Mag}) \times \mathcal{Q}(\mathbb{K}\mathbf{Mag}) \rightarrow \mathcal{Q}(\mathbb{K}\mathbf{Mag})$ defined by
 - ▷ $\theta_1 \wedge_i \theta_2 = \mathbb{K}\mathbf{Mag}/_{h_1+h_2}$;
 - ▷ $\theta_1 \vee_i \theta_2 = \mathbb{K}\mathbf{Mag}/_{h_1 \cap h_2}$.

Theorem [C.-Cordero-Giraudo, 2018]. Consider the notations introduced above.

- i. The tuple $(\mathcal{Q}(\mathbb{K}\mathbf{Mag}), \preceq_i, \wedge_i, \vee_i)$ is a lattice.
- ii. We have the following Grassmann formula analog:

$$\mathcal{H}_{\theta_1 \vee_i \theta_2}(t) + \mathcal{H}_{\theta_1 \wedge_i \theta_2}(t) = \mathcal{H}_{\theta_1}(t) + \mathcal{H}_{\theta_2}(t).$$

► Let $\mathbf{As} := \mathbb{K}\mathbf{Mag}/I_{\mathbf{As}}$ and $\mathbf{AAs} := \mathbb{K}\mathbf{Mag}/I_{\mathbf{AAs}}$, where $I_{\mathbf{As}}$ and $I_{\mathbf{AAs}}$ are generated by



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$$\begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} - \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} \quad \text{and} \quad \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} + \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array}$$

- Let $2\mathbf{Nil} := \mathbf{As} \wedge_i \mathbf{AAs}$, that is $I_{2\mathbf{Nil}} = I_{\mathbf{As}} + I_{\mathbf{AAs}}$;

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- ▶ using the Grassmann formula, one shows that π is an isomorphism.

III. Comb associative operads

► $\gamma \geq 1$: a positive integer;

► $I_{\text{CAs}(\gamma)}$: the ideal generated by



► $\gamma \geq 1$: a positive integer;

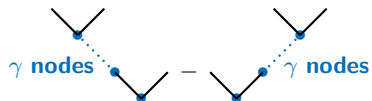
► $I_{\mathbf{CAs}(\gamma)}$: the ideal generated by



► $\mathbf{CAs}^{(\gamma)} := \mathbf{Mag}/I_{\mathbf{CAs}(\gamma)}$ is called the γ -comb associative operad.

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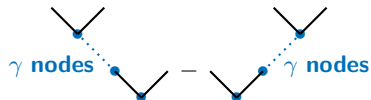
▷ $\mathbf{CAs}(\gamma) := \mathbf{Mag}/I_{\mathbf{CAs}(\gamma)}$ is called the γ -comb associative operad.

► For instance,

▷ $\mathbf{CAs}^{(1)} = \mathbb{K}\mathbf{Mag}$, $\mathbf{CAs}^{(2)} = \mathbf{As}$

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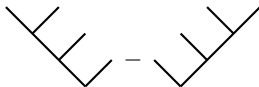
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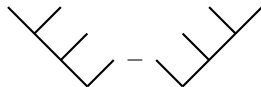
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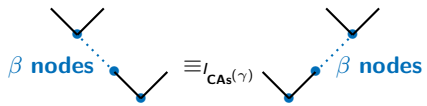
► Objective of the section: show that

$$\mathbf{CAs} := \{ \mathbf{CAs}^{(\gamma)} \mid \gamma \geq 1 \}$$

admits a lattice structure.

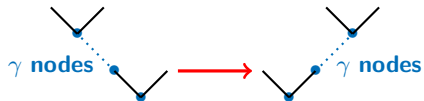
► \mathcal{A}_d : the restriction of \mathcal{A}_i to **CAs**;

▷ $\mathbf{CAs}^{(\gamma)} \mathcal{A}_d \mathbf{CAs}^{(\beta)}$ is equivalent to



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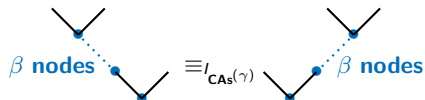
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▷ using an orientation of $\equiv_{\mathbf{CAs}^{(\gamma)}}$:

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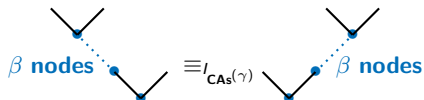
▷ **CAs** $^{(\gamma)}$ $\underline{\mathcal{A}}_d$ **CAs** $^{(\beta)}$ is equivalent to



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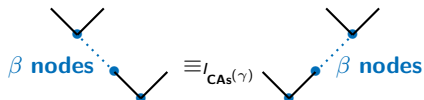
► Let $\wedge_d, \vee_d : \mathbf{CAs} \times \mathbf{CAs} \rightarrow \mathbf{CAs}$ defined by

▷ **CAs** $^{(\gamma)} \wedge_d$ **CAs** $^{(\beta)} := \mathbf{CAs}^{\left(\gcd(\bar{\gamma}, \bar{\beta}) + 1\right)}$;

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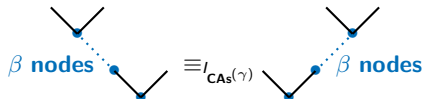
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Theorem [C.-Cordero-Giraud, 2018]. *The tuple $(\mathbf{CAs}, \preceq_d, \wedge_d, \vee_d)$ is a lattice.*

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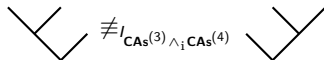
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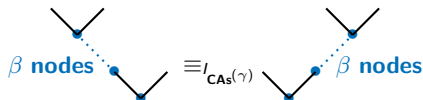
Theorem [C.-Cordero-Giraud, 2018]. *The tuple $(\mathbf{CAs}, \preceq_d, \wedge_d, \vee_d)$ is a lattice.*

Remark. $(\mathbf{CAs}, \preceq_d, \wedge_d, \vee_d)$ **does not embed** into $(\mathcal{Q}(\mathbb{K}\text{Mag}), \preceq_i, \wedge_i, \vee_i)$ as a sublattice:



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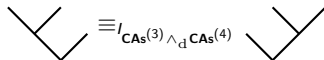
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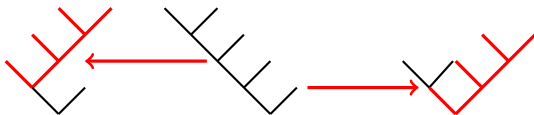
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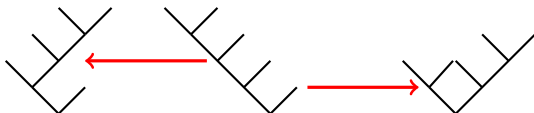


since $\mathbf{CAs}^{(3)} \wedge_d \mathbf{CAs}^{(4)} = \mathbf{CAs}^{(\gcd(2,3)+1)} = \mathbf{CAs}^{(2)} = \mathbf{As}$.

- The orientation of $\equiv_{I_{\text{CAs}}(\gamma)}$ is not confluent:

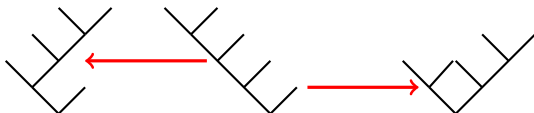


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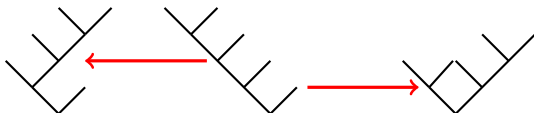
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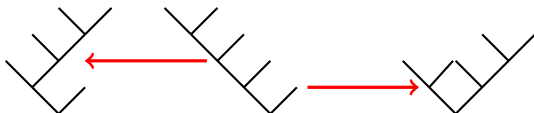
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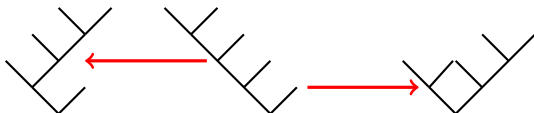
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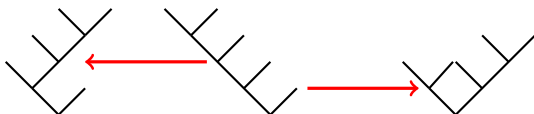
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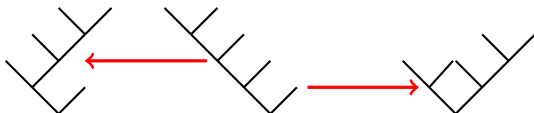
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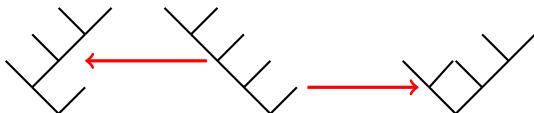
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IV. Conclusion and perspectives

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THANK YOU FOR LISTENING!