

# Second-order kernel online convex optimization with adaptive sketching

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## Motivation

- Non-parametric models are versatile and accurate
- First-order methods are fast but high regret
- Second-order methods suffer low regret but slow

$$\mathcal{O}(t^3) \text{ time } \mathcal{O}(t^2) \text{ space } (t \text{ steps})$$

- Current limitation: No interpretation for non-parametric regret, no approximate second-order methods

We propose **Sketched-KONS**, the first approximate algorithm for second-order Kernel Online Convex Optimization

- approximation  $\Rightarrow 1/\gamma$  times more regret but a  $\gamma^2$  speedup
- using a novel kernel matrix sketching technique
- regret scales with the effective dimension of the problem

## Kernel Online Convex Optimization

Online game between learner and adversary, at each round  $t \in [T]$

- the adversary reveals a new point  $\varphi(\mathbf{x}_t) = \phi_t \in \mathcal{H}$
- the learner chooses  $\mathbf{w}_t$  and predicts  $f_{\mathbf{w}_t}(\mathbf{x}_t) = \varphi(\mathbf{x}_t)^\top \mathbf{w}_t$ ,
- the adversary reveals the curved loss  $\ell_t$ ,
- the learner suffers  $\ell_t(\phi_t^\top \mathbf{w}_t)$  and observes gradient  $\mathbf{g}_t$ .

Kernel

- $\varphi(\cdot) : \mathcal{X} \rightarrow \mathcal{H}$  is the high-dimensional (possibly infinite) map
- $\Phi_t = [\phi_1, \dots, \phi_t]$ ,  $\Phi_t^\top \Phi_t = \mathbf{K}_t$  (kernel trick)
- $\mathbf{g}_t = \ell'_t(\phi_t^\top \mathbf{w}_t) \phi_t := \hat{g}_t \phi_t$

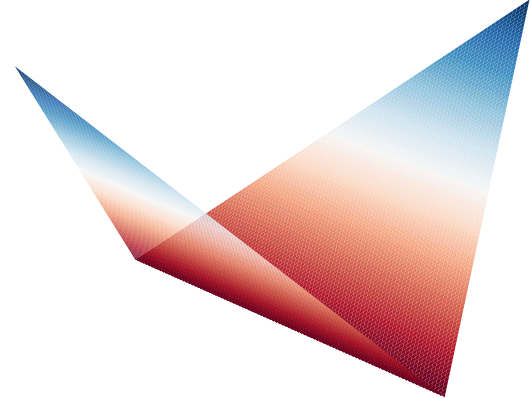
Minimize regret

$$R(\mathbf{w}) = \sum_{t=1}^T \ell_t(\phi_t^\top \mathbf{w}_t) - \ell_t(\phi_t^\top \mathbf{w})$$

against the best-in-hindsight  $\mathbf{w}^* := \arg \min_{\mathbf{w} \in \mathcal{H}} \sum_{t=1}^T \ell_t(\phi_t^\top \mathbf{w})$

## Curvature and first vs second order

Convex



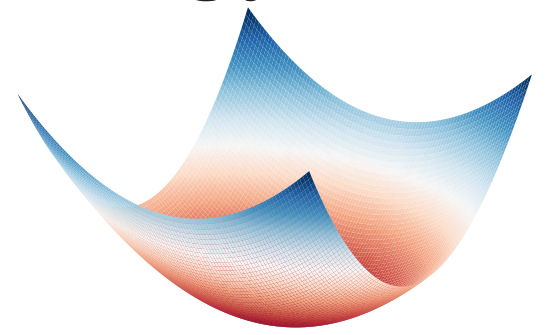
First order (GD)

Zinkevich 2003, Kivinen et al. 2004

- $\mathcal{O}(d)/\mathcal{O}(t)$  time/space per-step
- regret  $\sqrt{T}$

Approximation avoids  $\mathcal{O}(t)$  runtime  
 $\hookrightarrow$  but introduces approximation error (potentially  $\mathcal{O}(T)$  regret)

Strongly Convex



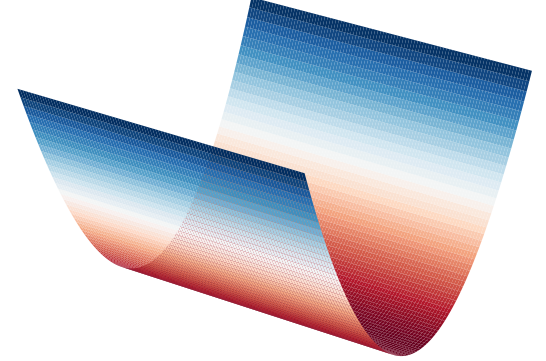
First order (GD)

Hazan, Rakhlin, et al. 2008

- $\mathcal{O}(d)/\mathcal{O}(t)$  time/space per-step
- regret  $\log(T)$

but often not satisfied in practice  
 $\hookrightarrow$  (e.g.  $(y_t - \phi_t^\top \mathbf{w}_t)^2$ )

$\sigma$ -curved



First order (GD)

- $\mathcal{O}(d)/\mathcal{O}(t)$  time/space per-step
- regret  $\sqrt{T}$

Second order (Newton-like)

Hazan, Kalai, et al. 2006, Zhdanov and Kalnishkan 2010

- regret  $\log(T)$
- $\mathcal{O}(d^2)/\mathcal{O}(t^2)$  time/space per-step

Fast approximations for linear case

Luo et al. 2016

$\hookrightarrow$  no approximate methods for kernel case

Assumptions

- the losses  $\ell_t$  are scalar Lipschitz  $|\ell'_t(z)| \leq L$
- $\ell_t(\phi_t^\top \mathbf{w}) \geq \ell_t(\phi_t^\top \mathbf{u}) + \nabla \ell_t(\phi_t^\top \mathbf{u})^\top (\mathbf{w} - \mathbf{u}) + \sigma (\nabla \ell_t(\phi_t^\top \mathbf{u})^\top (\mathbf{w} - \mathbf{u}))^2$

Challenge

Reduce computational cost without losing logarithmic regret?

## References

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- Haipeng Luo et al. "Efficient second-order online learning via sketching". In: Neural Information Processing Systems. 2016.
- Fedor Zhdanov and Yuri Kalnishkan. "An Identity for Kernel Ridge Regression". In: Algorithmic Learning Theory. 2010.
- Martin Zinkevich. "Online Convex Programming and Generalized Infinitesimal Gradient Ascent". In: ICML. 2003.

## Kernel Online Newton Step (KONS)

Second-Order Gradient Descent

- $\mathbf{A}_0 = \alpha \mathbf{I}$
- $\mathbf{A}_t = \mathbf{A}_{t-1} + \sigma \mathbf{g}_t \mathbf{g}_t^\top$
- $\mathbf{w}_{t+1} = \mathbf{w}_t - \mathbf{A}_t^{-1} \mathbf{g}_t$

$$\mathbf{A}_t^{-1} = \left( \begin{array}{c|c} \mathbf{A}_{t-1}^{-1} & \\ \hline & \sigma \mathbf{g}_t \mathbf{g}_t^\top \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \right)^{-1} = \begin{array}{c} \mathbf{A}_{t-1}^{-1} \\ \hline \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \left( \begin{array}{c} \mathbf{A}_{t-1}^{-1} \\ \hline \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \right)^{-1} \begin{array}{c} \mathbf{I} \\ \hline \end{array}$$

$$R(\mathbf{w}) \leq \mathcal{O} \left( \sum_{t=1}^T \mathbf{g}_t^\top \mathbf{A}_t^{-1} \mathbf{g}_t \right) \leq \mathcal{O} \left( \sum_{t=1}^T \mathbf{g}_t^\top (\mathbf{G}_t \mathbf{G}_t^\top + \alpha \mathbf{I})^{-1} \mathbf{g}_t \right) \leq \mathcal{O} \left( L \sum_{t=1}^T \phi_t^\top (\Phi_t \Phi_t^\top + \alpha \mathbf{I})^{-1} \phi_t \right) \leq \begin{cases} \text{LOCO: } \mathcal{O}(d \log(T)) \\ \text{KOCO: } \mathcal{O}(\log(\text{Det}(\mathbf{K}_T + \alpha \mathbf{I}))) \end{cases}$$

## Effective dimension

Lemma 1

$$d_{\text{onl}}^T(\alpha) := \sum_{t=1}^T \phi_t^\top (\Phi_t \Phi_t^\top + \alpha \mathbf{I})^{-1} \phi_t \leq \log(\text{Det}(\mathbf{K}_T / \alpha + \mathbf{I})) \leq 2d_{\text{eff}}^T(\alpha) \log(T/\alpha).$$

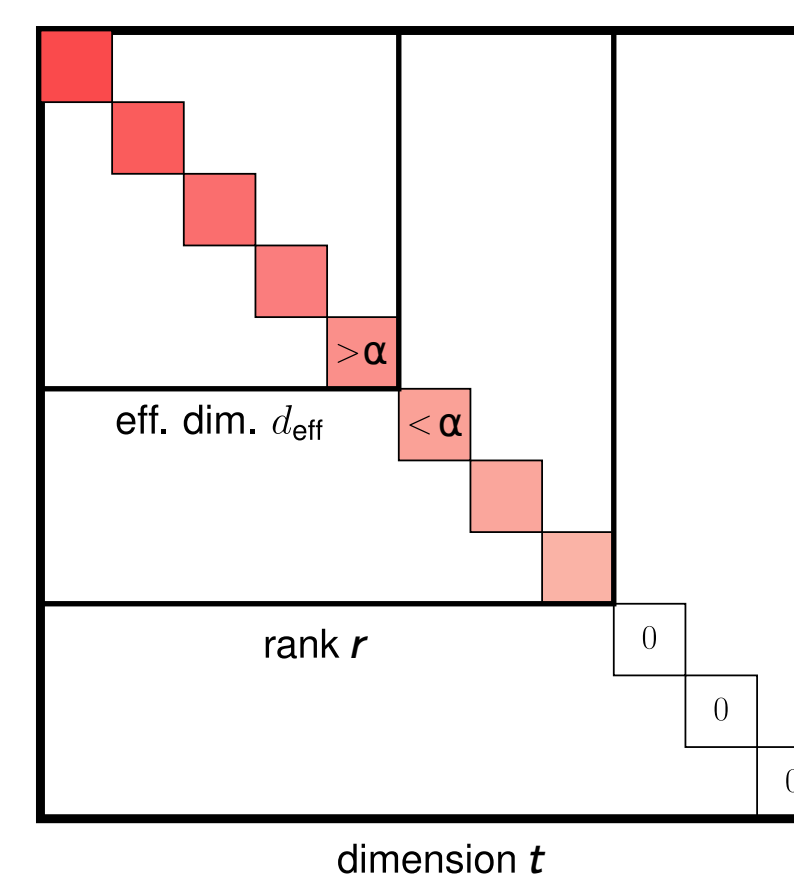
Given a kernel matrix  $\mathbf{K}_T \in \mathbb{R}^{t \times t}$

$\Rightarrow$   $\alpha$ -ridge leverage score

$$\tau_{T,i}(\alpha) = \mathbf{e}_{T,i}^\top \mathbf{K}_T (\mathbf{K}_T + \alpha \mathbf{I}_T)^{-1} \mathbf{e}_{T,i} = \phi_i^\top (\Phi_T \Phi_T^\top + \alpha \mathbf{I})^{-1} \phi_i$$

$\Rightarrow$  Effective dimension

$$d_{\text{eff}}(\alpha)_T = \sum_{i=1}^T \tau_{T,i}(\alpha) = \text{Tr}(\mathbf{K}_T (\mathbf{K}_T + \alpha \mathbf{I}_T)^{-1}) = \sum_{i=1}^T \frac{\lambda_i(\mathbf{K}_T)}{\lambda_i(\mathbf{K}_T) + \alpha} \leq \text{Rank}(\mathbf{K}_T) = r$$



## Kernel Online Row Sampling (KORS)

Input: Regularization  $\alpha$ , accuracy  $\varepsilon$ , budget  $\beta$

- Initialize  $\mathcal{I}_0 = \emptyset$
- for  $t = \{0, \dots, T-1\}$  do
- receive  $\tilde{\phi}_t$
- construct temporary dictionary  $\tilde{\mathcal{I}}_t := \mathcal{I}_{t-1} \cup (t, 1)$
- compute  $\tilde{p}_t = \min\{\beta \tilde{\tau}_{t,t}, 1\}$  using  $\tilde{\mathcal{I}}_t$  and Eq. 4 in the paper.
- draw  $z_t \sim \mathcal{B}(\tilde{p}_t)$  and if  $z_t = 1$ , add  $(t, 1/\tilde{p}_t)$  to  $\mathcal{I}_t$
- end for

Theorem 1. Given parameters  $0 < \varepsilon \leq 1$ ,  $0 < \alpha$ ,  $0 < \delta < 1$ , let  $\rho = \frac{1+\varepsilon}{1-\varepsilon}$  and run KORS with  $\beta \geq 3 \log(T/\delta)/\varepsilon^2$ . Then w.p.  $1 - \delta$ , for all steps  $t \in [T]$ ,

- $(1 - \varepsilon) \mathbf{A}_t \preceq \mathbf{A}_t^T \preceq (1 + \varepsilon) \mathbf{A}_t$ .
- $|\mathcal{I}_t| \leq d_{\text{eff}}^T(\alpha) \frac{6\rho \log^2(\frac{2T}{\delta})}{\varepsilon^2}$ .
- Satisfies  $\tau_{t,t} \leq \tilde{\tau}_{t,t} \leq \rho \tau_{t,t}$ .

Moreover, the algorithm runs in  $\mathcal{O}(d_{\text{eff}}^T(\alpha)^2 \log^4(T))$  space, and  $\mathcal{O}(d_{\text{eff}}^T(\alpha)^2 \log^4(T))$  time per iteration.

## Sketched-KONS

Naive Approach:  $\tilde{\mathbf{A}}_t = \tilde{\mathbf{A}}_{t-1} + (\mathbb{I}\{\text{coin flip w.p. } p_t\} / p_t) \sigma \mathbf{g}_t \mathbf{g}_t^\top$  with  $p_t \propto \tilde{\tau}_{t,t}$

$$\tilde{\mathbf{A}}_t^{-1} = \left( \begin{array}{c|c} \tilde{\mathbf{A}}_{t-1}^{-1} & \\ \hline & \sigma \mathbf{g}_t \mathbf{g}_t^\top / p_t \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \right)^{-1} = \begin{array}{c} \tilde{\mathbf{A}}_{t-1}^{-1} \\ \hline \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \left( \begin{array}{c} \tilde{\mathbf{A}}_{t-1}^{-1} \\ \hline \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \right)^{-1} \begin{array}{c} \mathbf{I} \\ \hline \end{array}$$

- w.h.p.  $\tilde{\mathbf{A}}_t$  updated only  $d_{\text{eff}}^T(\alpha) \log^2(T)$  times
- $\tilde{\mathcal{O}}(d_{\text{eff}}^T(\alpha)^2 + t)$  per-step space/time complexity

- Expected regret  $d_{\text{eff}}^T(\alpha) \log(T)$
- The weights  $1/p_t \sim 1/\tilde{\tau}_{t,t}$  can be large  
 $\hookrightarrow$  large variance

SKETCHED-KONS  $\tilde{\mathbf{A}}_t = \tilde{\mathbf{A}}_{t-1} + (\mathbb{I}\{\text{coin flip w.p. } p_t\} / p_t) \sigma \mathbf{g}_t \mathbf{g}_t^\top$  with  $p_t \propto \max\{\gamma, \tilde{\tau}_{t,t}\}$

$$\tilde{\mathbf{A}}_t^{-1} = \left( \begin{array}{c|c} \tilde{\mathbf{A}}_{t-1}^{-1} & \\ \hline & \sigma \mathbf{g}_t \mathbf{g}_t^\top / p_t \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \right)^{-1} = \begin{array}{c} \tilde{\mathbf{A}}_{t-1}^{-1} \\ \hline \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \left( \begin{array}{c} \tilde{\mathbf{A}}_{t-1}^{-1} \\ \hline \end{array} + \begin{array}{c} \mathbf{I} \\ \hline \end{array} \right)^{-1} \begin{array}{c} \mathbf{I} \\ \hline \end{array}$$

Theorem 2. For any sequence of losses  $\ell_t$  satisfying Asm. 1-2, let  $\tilde{\tau}_{\min} = \min_{t=1}^T \tilde{\tau}_{t,t}$ . For all  $t$ ,  $\alpha \leq \sqrt{T}$ ,  $\beta \geq 3 \log(T/\delta)/\varepsilon^2$ , then w.p.  $1 - \delta$  the regret of SKETCHED-KONS satisfies

$$\tilde{R}_T \leq \alpha \|\mathbf{w}^*\|^2 + 2 \frac{d_{\text{eff}}^T(\alpha / (\sigma L^2)) \log(2\sigma L^2 T)}{\sigma \max\{\gamma, \beta \tilde{\tau}_{\min}\}}, \quad (1)$$

and the algorithm runs in  $\mathcal{O}(d_{\text{eff}}^T(\alpha)^2 + t^2 \gamma^2)$  time and  $\mathcal{O}(d_{\text{eff}}^T(\alpha)^2 + t^2 \gamma^2)$  space complexity for each iteration  $t$ .

- Trade-off computation and regret  
 $\hookrightarrow 1/\gamma$  increase in regret for  $\gamma^2$  space/time improvement
- Neither uniform nor RLS  
 $\hookrightarrow$  keep updates with high  $\tau_{t,t}$  for accuracy  
 uniformly update for stability
- Can we get rid of dependency on  $t$ ?  
 $\hookrightarrow$  not when  $\mathbf{A}_t - \mathbf{A}_{t-1} = w_t \mathbf{g}_t \mathbf{g}_t^\top$

## Counterexample

Adversary always plays same sample  $\phi_{\text{exp}}$ , but alternates label  $\{+1, -1\}$

Class of updates:  $\mathbf{A}_t - \mathbf{A}_{t-1} = w_t \mathbf{g}_t$

#SV budget  $B = \#\mathbb{I}\{w_t \neq 0\}$  drives complexity  
 cumulative weight  $W_t = \sum_{s=1}^t w_s$  drives regret

$$R(\mathbf{w}^*) \leq \sum_{t=1}^T \mathbf{g}_t^\top \mathbf{A}_t^{-1} \mathbf{g}_t + \sum_{t=1}^T (\mathbf{w}_t - \mathbf{w}^*)^\top (\mathbf{A}_t - \mathbf{A}_{t-1} - \sigma_t \mathbf{g}_t \mathbf{g}_t^\top) (\mathbf{w}_t - \mathbf{w}^*) \leq \sum_{t=1}^T \mathbf{g}_t^\top \mathbf{A}_t^{-1} \mathbf{g}_t + \sum_{t=1}^T (w_t - \sigma_t)^2 (\mathbf{g}_t^\top (\mathbf{w}_t - \mathbf{w}^*))^2 \leq \underbrace{\sum_{s=1}^t \frac{1}{W_s + \alpha}}_{R_G} + \underbrace{\sum_{s=1}^t \max\{0, w_t - \sigma\}}_{R_D}$$

- Increase  $W_t$  quickly  
 $\hookrightarrow$  reduce  $R_G$
- Increase  $W_t$  slowly  
 $\hookrightarrow$  reduce  $R_D$
- Increase  $W_t$  sparsely  
 $\hookrightarrow$  reduce  $B$

Contrasting goals cannot be satisfied at the same time.

Only constant speedup over exact

HOW CAN WE AVOID THIS?

Support Removal

Learn how to remove old  $\mathbf{g}_{t-1}$  from  $\mathbf{A}_t$ ?  
 $\hookrightarrow (\mathbf{w}_t - \mathbf{w}^*)^\top (\mathbf{g}_t \mathbf{g}_t^\top - \mathbf{g}_{t-1} \mathbf{g}_{t-1}^\top) (\mathbf{w}_t - \mathbf{w}^*)$  could be large

Functional embedding

Instead of approximating  $\mathbf{A}_t$ , approximate  $\phi_t$   
 $\hookrightarrow$  Random features not strong enough (yet)

Avron et al. ICML'17 satisfy guarantee (1) of Thm. 1  
 $\hookrightarrow$  only in batch setting

Nyström-based embeddings?

$\hookrightarrow$  ongoing work